Pattern-based Rewrite and Refinement of Architectures Using Graph Theory



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Introduction & Motivation



3.6.2 Message-oriented Middleware

Asynchronous message-based communication is provided while hiding complexity resulting from addressing, routing, or data formats from communication partners to make interaction robust and flexible.

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How can communication partners exchange information asynchronously with a communication partner?

Context

The application components of a distributed application (139) are hosted on multiple cloud resources and have to exchange information with each other. Often, the integration with other cloud applications and non-cloud applications is also required. These different applications possibly use different programming languages, data formats, and execution environments (94). When one application directly exchanges information with another application the address and data format of the target application has to be respected. Even within one homogeneous distributed application (139), the communicating application components must be available at the time when the information shall be exchanged. These dependencies can significantly reduce the availability of the overall application as the failure of one application component would directly affect all components communications with it. The resulting dependency between communication partners regarding their location, availability, and data format is called *tight coupling*. It also increases the complexity of the management of the overall application or the landscape of applications, because changes to one communication partner, for example, regarding the format of exchanged data or the address used, also affect the other communication partner directly. This should be avoided to increase the availability of the overall application and ease continuous alterations by making communication flexible

Solution

Communication partners exchange information asynchronously using messages handled by a message originated middleware. For this purpose a message originated

A pub-sub channel [1] may be used to broadcast a message to multiple receivers. While a message queue conceptually delivers messages to only one receiver, a pub-sub channel delivers messages to multiple receivers.



Figure 37: Message-oriented Middleware and Related Patterns

Result

When interacting with a message-oriented middleware, a sender puts a message on one message queue or pub-sub channel and receivers can retrieve it from possibly different queues. In between these two access points the message-oriented middleware handles the complexity of addressing, availability of communication partners and message format transformation as shown in Figure 37. Therefore, in addition to message queues, the message-oriented middleware provides components that route messages to intended receivers as well as handle message format transformation. Communication partners may communicate via messages without the need to know the message format expected by the communication partner or the address at which it can be reached. Furthermore, communication partners can send and receive messages at their own pace and without relying on the availability of communication partners.

[1]The message-oriented middleware, therefore, suggests a pipes-and-filters application architecture as covered by Hohpe and Woolf [1], Bushmann et al. [14] and in the distributed application (143) pattern in Chapter 4. In this scope, application components act as independently operating *filters* that are interconnected through *pipes*, i.e., the message queues provided by the *message-oriented middle-ware*. Hohpe and Woolf [1] describe the behavior of these pipes and how they may be connected. Regarding the *filters*, i.e., the application components, Hohpe and Woolf also described how to interface with the messaging systems in the *adapter* pattern.

The more intermediaries a message passes through while traversing a *message-oriented middleware*, the more likely it becomes that an intermediary fails. To address this issue, messages are often stored in persistent storage by *the message-oriented middleware* from where they can be recovered in case of failures. This approach is described by the *guaranteed delivery* pattern introduced by Hohpe and Woolf [1].

Variations

A message-oriented middleware is used if small amounts of data need to be exchanged frequently, as messages are often restricted in size so they can be handled more easily. If larger amounts of data have to be exchanged, messages may either contain a pointer to this data that is actually stored at a different location, for example, a storage offering (see Section 3.5) or the data may be split up among multiple messages. Hohpe and Wolf [1] cover patterns for this exchange of large data elements: the *file transfer* pattern describes how data may be exported from one application and imported by a different one. A message sequence may be used to split large data elements among a set of messages.

Related Patterns

Figure 37 shows how the other messaging patterns of this section are related to the message oriented middleware:

At-least-once delivery (128): it is ensured that messages traversing the message-oriented middleware are delivered once or multiple times. This is achieved through acknowledgements for message receives. If an acknowledgement is not received, a message is retransmitted.

Exactly-once delivery (126): messages traversing the message-oriented middleware are delivered once and only once to the receiver. This involves reliably storage of messages in the message-oriented middleware and, often, transactional message exchange during its traversal of the message-oriented middleware

In addition to these communication pattern should be considered to be implemented in teracting with a message-oriented middlew — Transaction-based processor (179): if

- Transaction-based processor (179): if transactions to assure that messages transaction can be extended to include the receiver as well. Therefore, the tr application to assure that messages ar havior is also described by Hohpe ar tern. The transaction-based processor extends it to the transactional interaction Timeout-based message processor (18 assures at-least-once delivery (128) b client can extend the acknowledgmen Therefore the timeout-based message sure that messages are processed at lea Distributed application (139): applic loosely coupled (139) application co oriented middleware to exchange inf scope, an idempotent processor (176 messages created by a message-orier deliverv
- Message mover component (201): thi different message-oriented middlewa providers or that are installed in on-pre Watchdog (230): a watchdog may be pecially, in scope of environment-ba queues to store information securely e Batch processing component (165): m delay messages. A batch processing c only when conditions are feasible, fo low or the overall application experien

Known Uses

Using messaging to integrate distributed approach. Many additional messaging patt

C. Fehling, F. Leymann, R. Retter, W. Schupeck, and P. Arbitter, Cloud Computing Patterns: Fundamentals to Design, Build, and Manage Cloud Applications. Springer, 2014.

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Application Example

Application Example – *Input*



Application Example – *M*₁: *MOM* – *Check Requirements*



Application Example – *M*₁: *MOM* – *Refine or Rewrite*



Application Example – M₂: Watchdog – Check Requirements



Application Example – M₂: Watchdog – Refine or Rewrite



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Concept & Formalization



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Conclusion & Future Work – *Poster Session*



Thank you for your attention – I hope for interesting discussions during the poster session!

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