

How to model the internet of everything

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How to model the internet of everything

yesterday:

Rick Kazman:

Challenges in Architecting

Web 4.0 systems

*“enormous design debts
in today’s software”*

*“How can we analyze this
kind of (software) systems?”*

Don’t do it!

Start with modeling

- the real world problem involved,
- as well as its solution!

*“make the company’s
administration more efficient”*

“How can we analyze models?”

Generate software from models!

“This has been attempted
several times
and always failed”

!Try harder!

Use the right
modeling concepts!

Which ones?

You will see

How to model the internet of everything

All engineering sciences
start with models!

CS starts with the product.

usual way:

First you write a lousy software,
then you try to repair it.

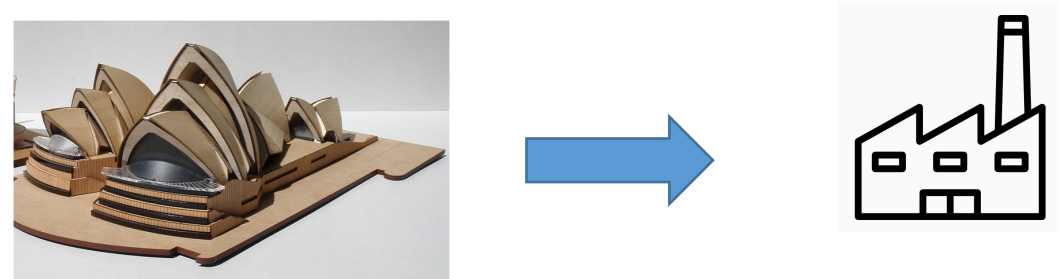
Modeling in informatics
is about changing the perspective.

Car production:



Can you imagine?

Instead:

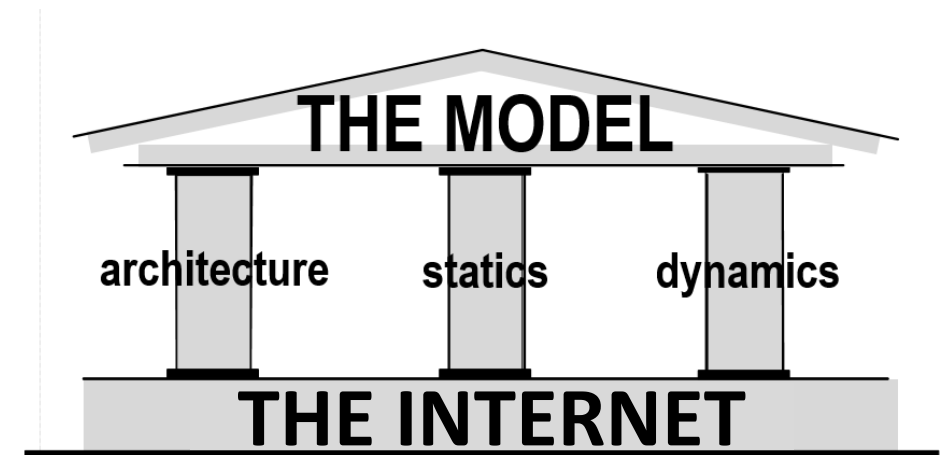


How to model the internet of everything

Three postulates:

a comprehensive model of the internet ...

- is *structured* **architecture**
- includes data ***and*** *real life items* **statics**
- is *locally* updated. **dynamics**



1. Architecture

modeling techniques

with modules and parallel composition

- process algebras,
- csp,
- statecharts

business process modeling

- EPC
- UML-AD
- RAD
- IDE

Architecture languages

- ABACUS
- ACME / ADML
- [ArchiMate](#)
- [ByADL](#)
- [DEMO](#)
- [LePUS3 and Class-Z](#)
- Rapide
- Wright

**all of them focus implementability,
not (big) systems!**

**we need a fundamental technique more general,
more abstract,
more simple**

focussing the real world problem

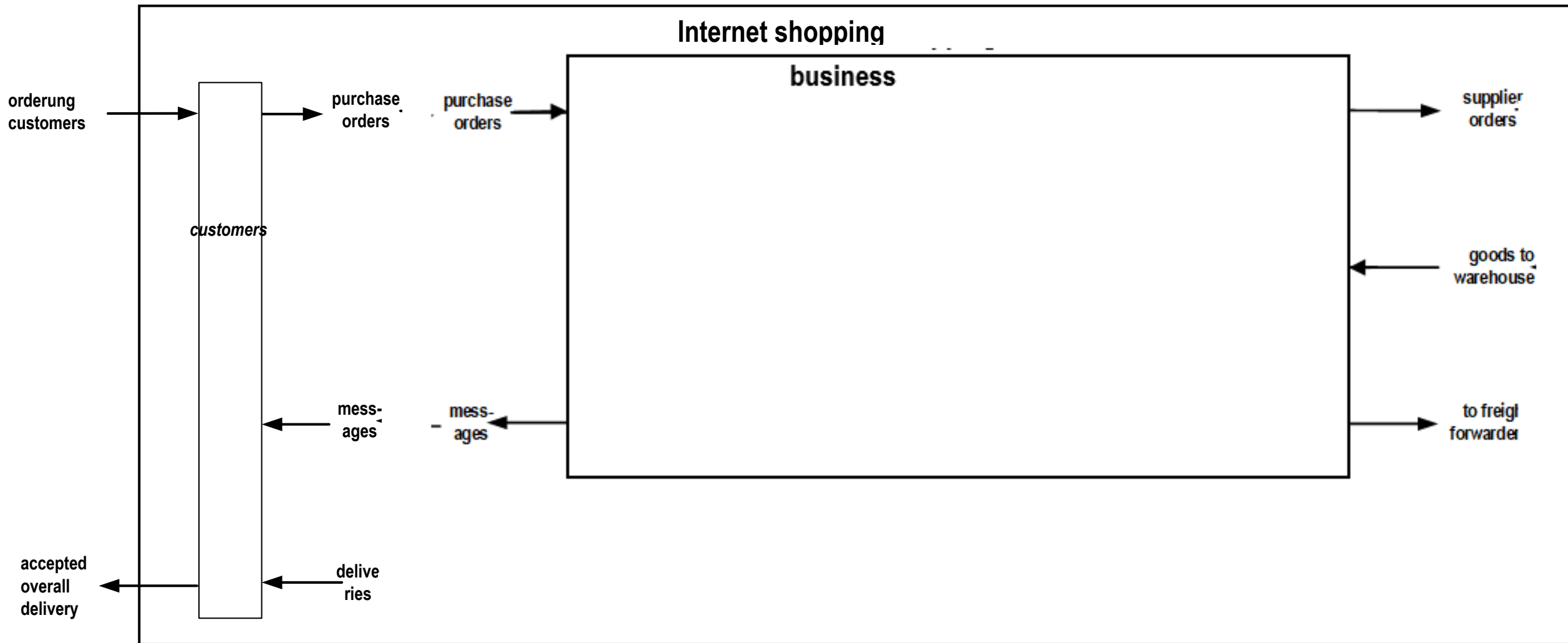
Internet shopping

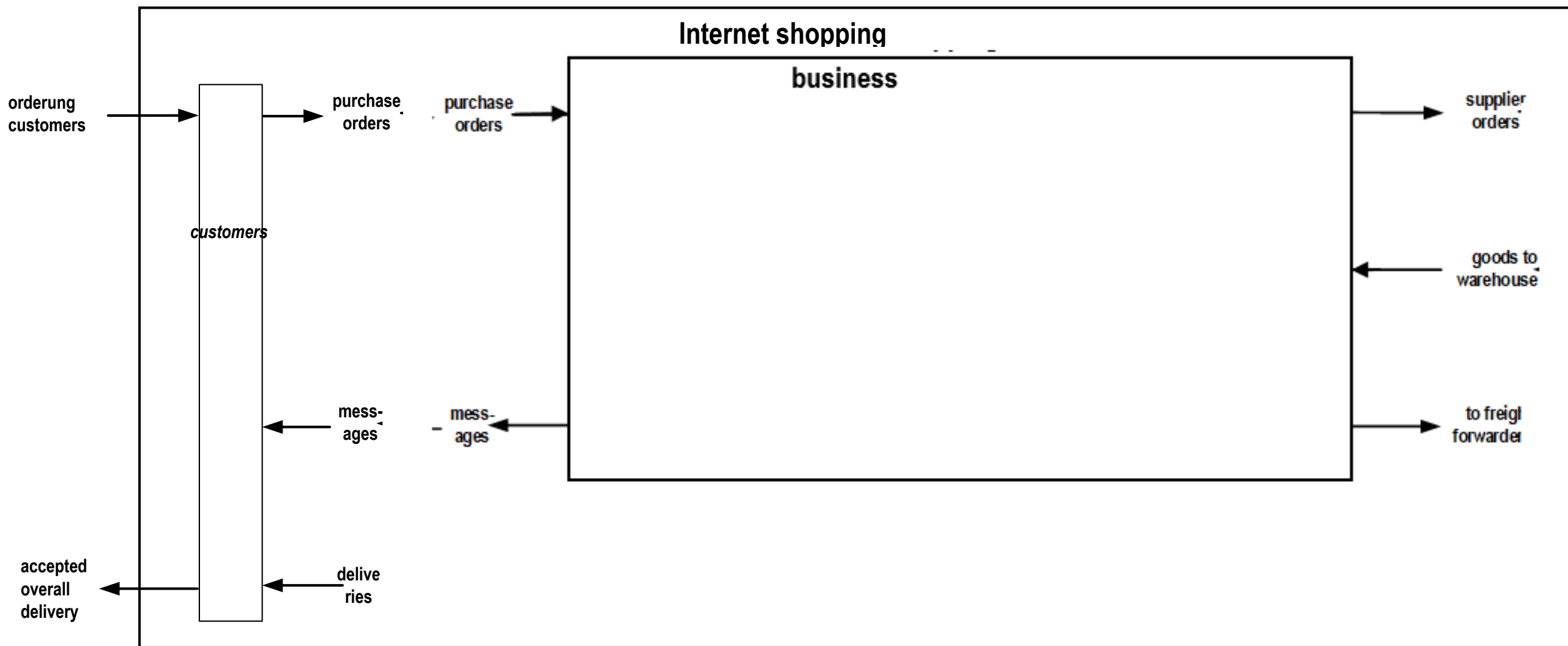
ordering
customers

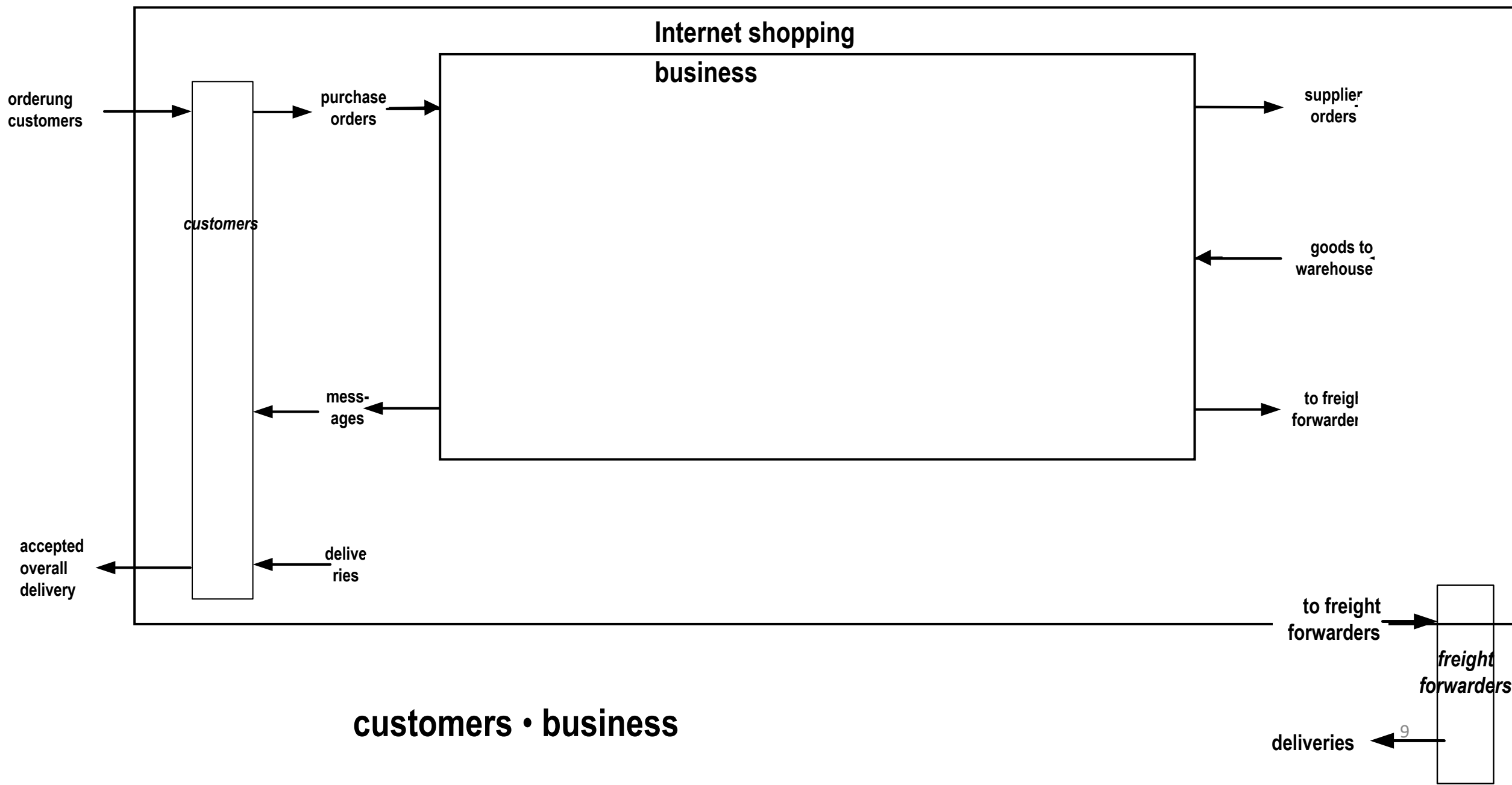


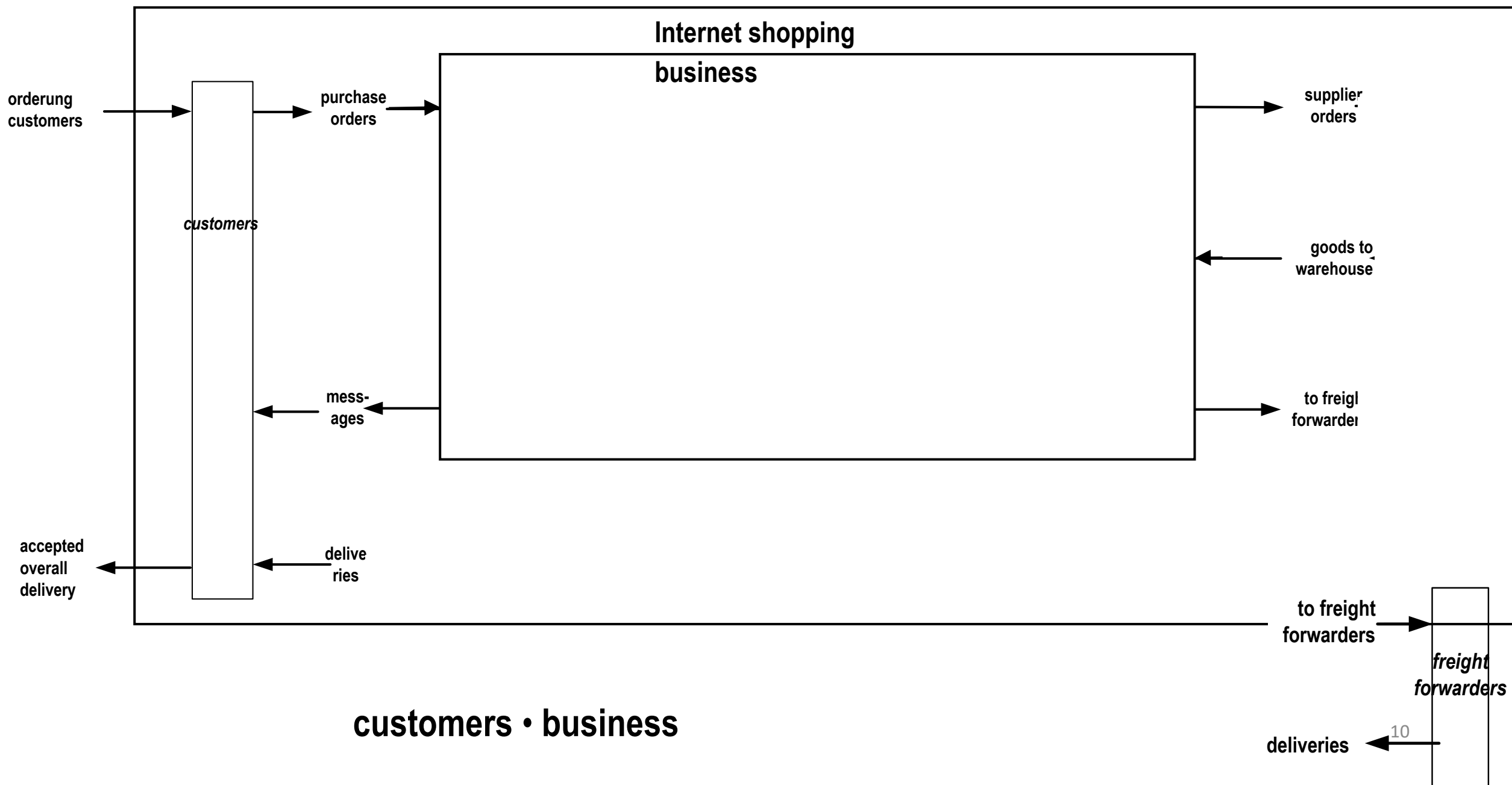
accepted
overall
delivery

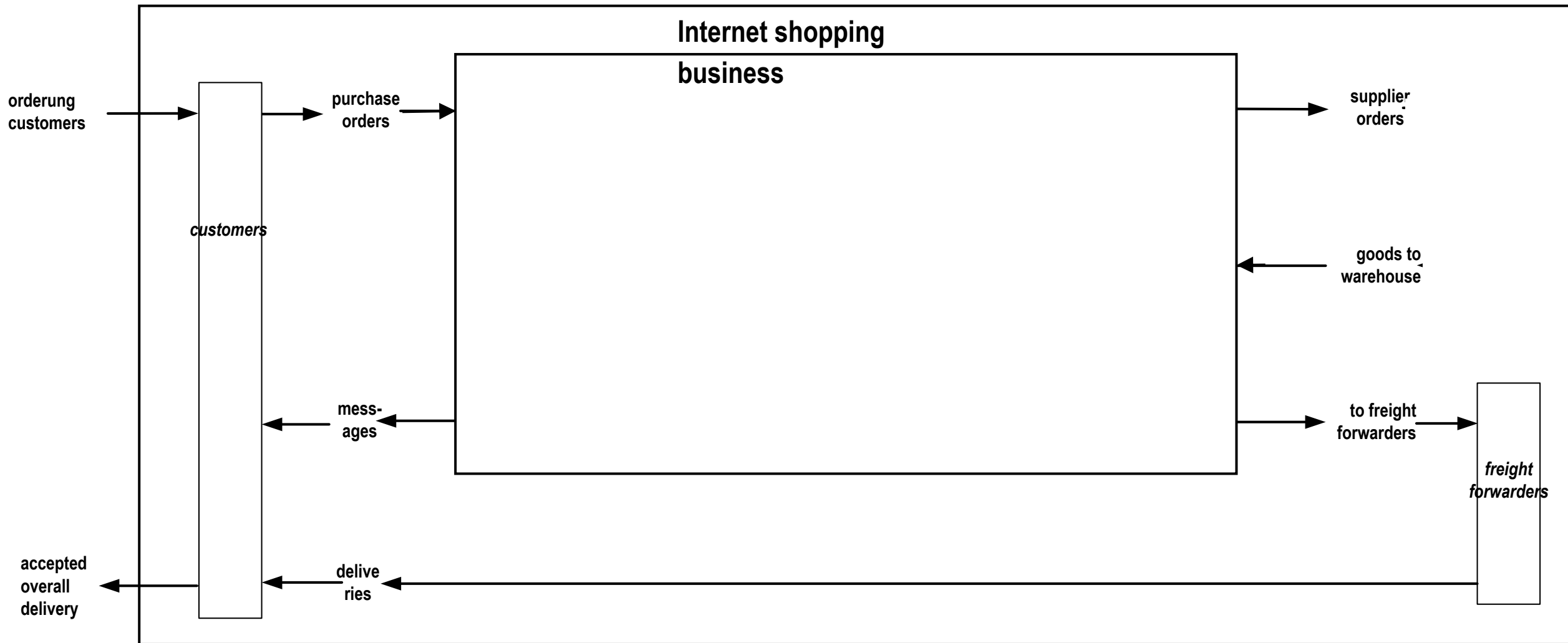




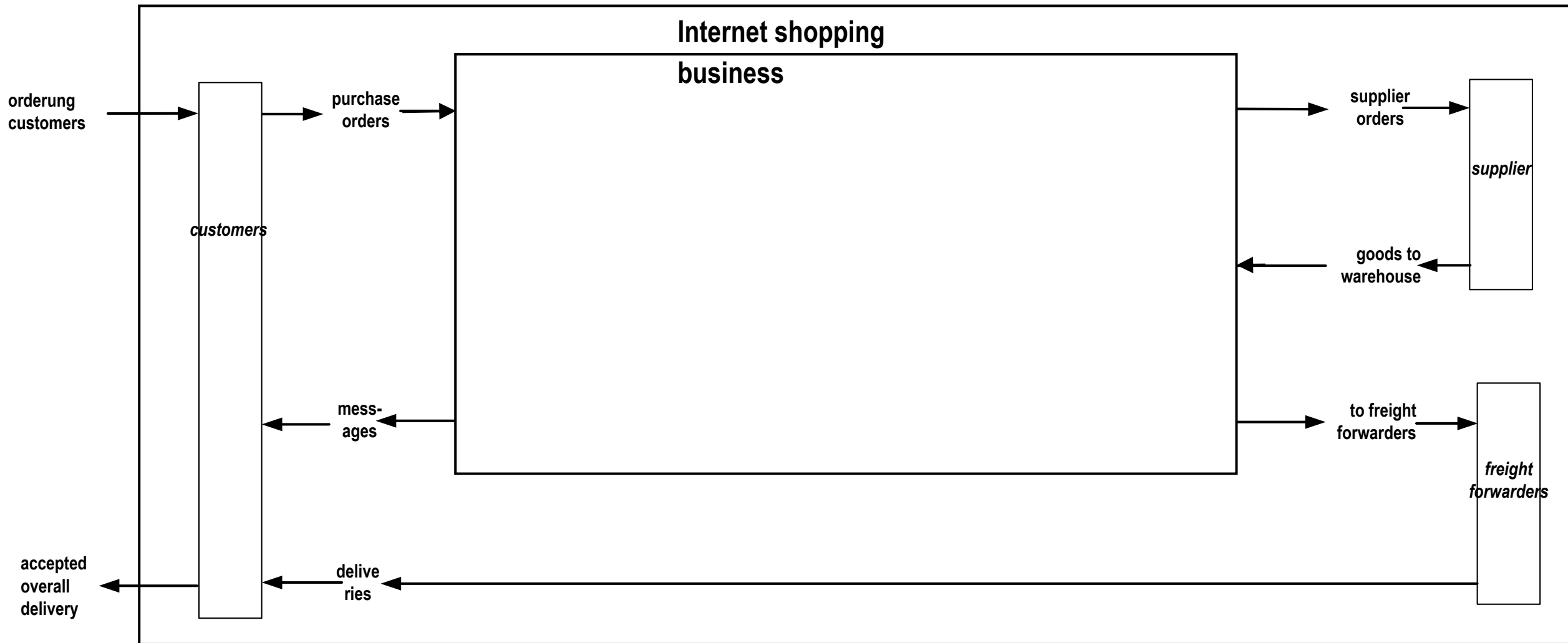




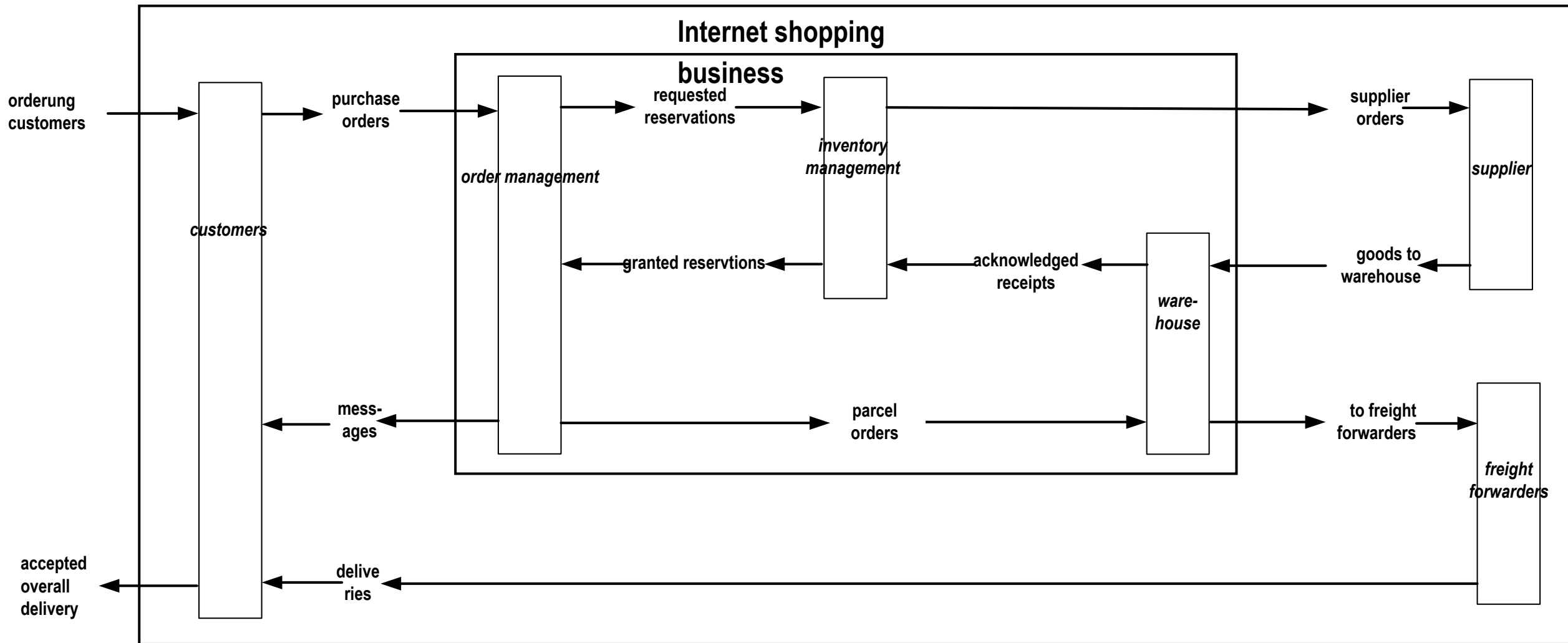




customers • business • freight forwarders



customers • business • freight forwarders • suppliers



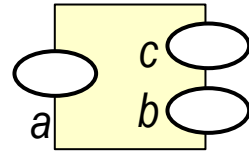
customers • business • freight forwarders • suppliers

where business = order management • inventory management • warehouse

1 Architecture: modules and composition

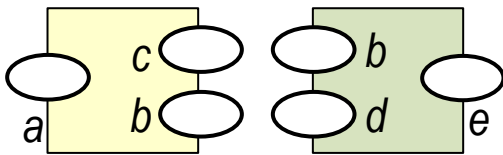
Def. Module:

graph G with two distinguished
sets of labelled nodes, *G and G^*
left and *right* interface



Theorem. $(G \bullet H) \bullet K = G \bullet (H \bullet K)$.

Def.: Composition $G \bullet H$ of modules G and H :
merge equally labelled nodes of G^* and *H .



1. Architecture: modules are intuitive!

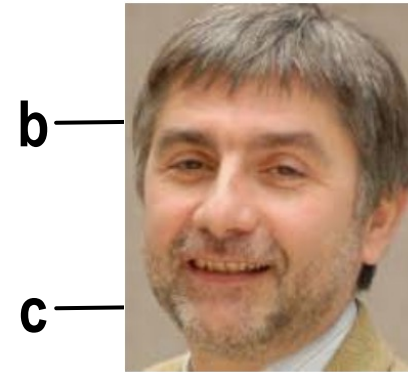
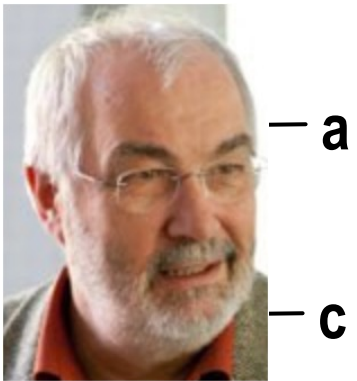
left and a *right* interface are quite natural:

G* and *G

- *input* and *output*,
- *customer* and *supplier*,
- *requester* and *provider*,
- *consumer* and *producer*,
- *buy side* and *sell side*,
- *predecessor* and *successor*,
- *assumptions* and *guarantees*,
- *pull* and *push*,
- *left* and *right*,

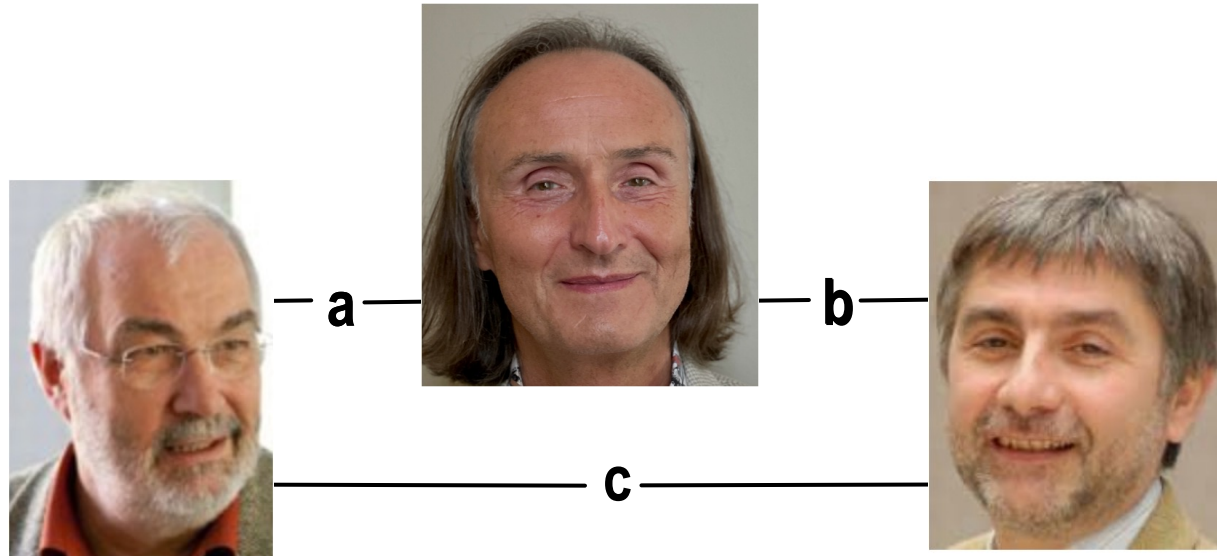
etc.

1. Architecture: Example



W • F • B

1. Architecture: Example



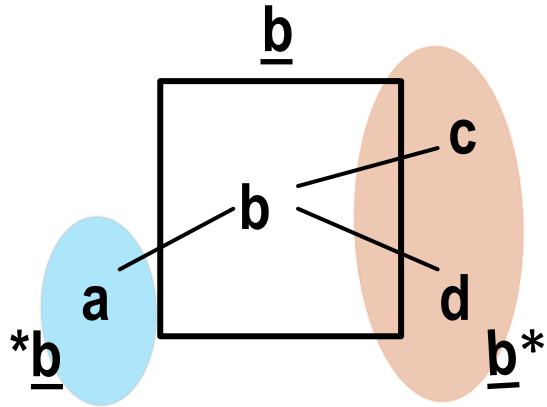
W • F • B



1. Architecture: composition is universal!

Each graph node yields node modules!

*a node module b of node **b***



Theorem. Each graph can be composed from node modules

$$\underline{a} \bullet \underline{b} \bullet \underline{c} \bullet \underline{d}$$

interfaces?

Composition with ...

- *one interface: not associative* $(G \bullet H) \bullet K \neq G \bullet (H \bullet K)$.
- *three interfaces: - not necessary (above Theorem)*
- technically intricate
- **two** interfaces: the basis
for *any* kind of composition!

1. Architecture: how manage complex composition?

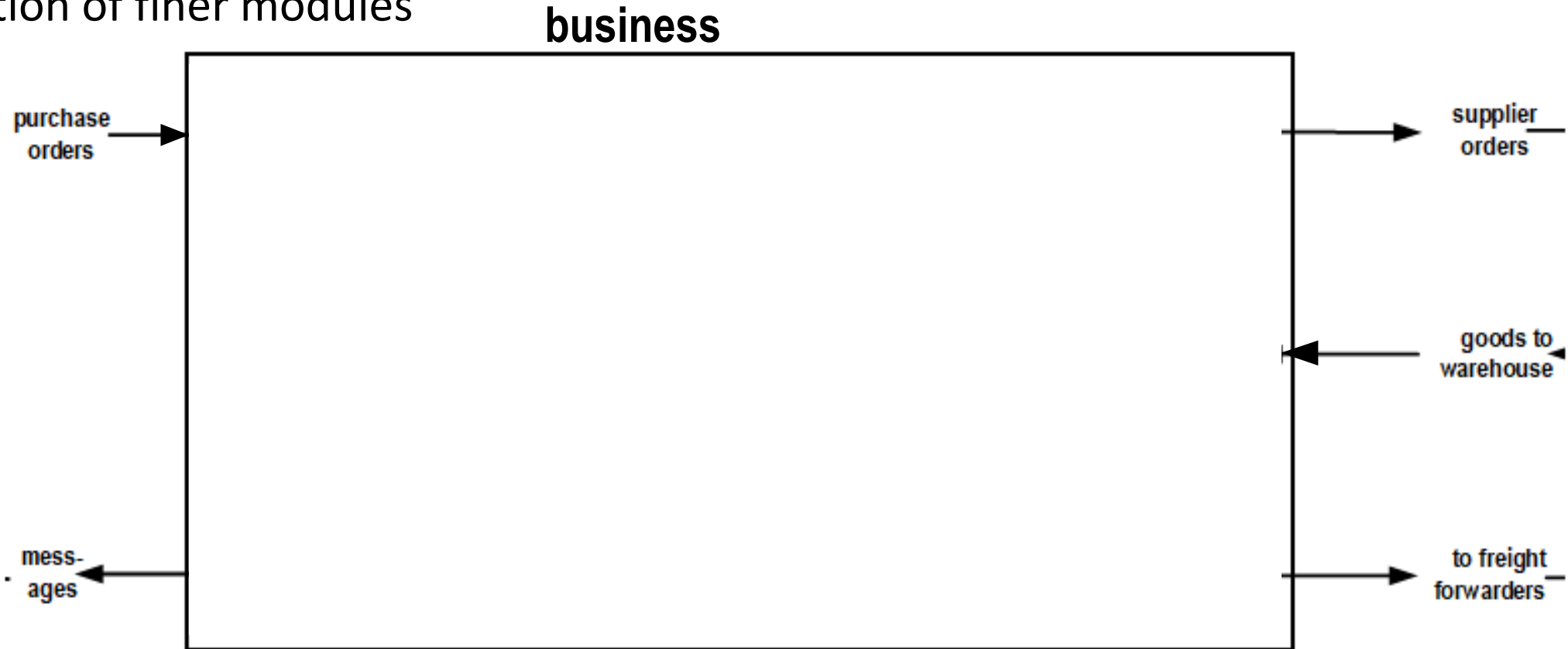
... by help of adapters:

replace $A * B$

by $A \bullet \text{adapter} \bullet B$,

1. Architecture: how manage refinement?

... by composition of finer modules



business = order management • inventory management • warehouse

1. Architecture: how manage “really big” systems?

by keeping every concept *local, inside modules!*

- level of abstraction
- composition
- refinement
- adapter
- data (few global names)
- behavior (no global states)

2. Statics

statically given:

- data,
- real life items.
- operations on data and items.

How organize all this?

Learn from Logic!

Structures and Signatures

a structure

ground sets

functions

derived sets

constants

another structure

ground sets

KN = all id card holders *customers*
AR = {small car, medium, large} *articles*
WA = {VW1, VW2, ... FORD1, FORD2, MERCEDES} *goods*
TE = {1.1. ... 31.12.} *dates*
SP = {Maier, Müller, Schulz} *freight forwarders*

derived sets

AP = **AR** x **N** *article items*
AL = **M**(**AP**) *article lists*
AM = **M**(**AR**) *sets of articles*
WM = **M**(**WA**) *sets of goods*

constants

p1: **AP** = (small car, 4),
p2: **AP** = (big, 2)
K: **P**(**KN**) = {Peter} *ordering customers*
G: **AL** = {(small car, 30), (medium, 2), (large, 1)} *initially listed articles*
H: **WM** = {VW, VW, BMW, BMW} *initially listed goods*
R: **P**(**SP**) = {Maier, Müller} *initally available freight forwarders*

functions

f(**w**) = „w“, for **w** ∈ **WA**

(a,n)' = **n**•[**a**] for (**a,n**) ∈ **AP**

a Σ -structure

ground sets

KN = all id card holders *customers*
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derived sets

AP = **AR** \times \mathbb{N} *article items*
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a signature, Σ

ground sorts

KN *customers*
AR *articles*
WA *goods*
TE *dates*
SP *freight forwarders*

derived sorts

AP = **AR** \times \mathbb{N} *article items*
AL = $M(\mathbf{AP})$ *article lists*
AM = $M(\mathbf{AR})$ *sets of articles*
WM = $M(\mathbf{WA})$ *sets of goods*

constant symbols

p1, p2: **AP** *ordered article positions*
B: $P(\mathbf{KN})$ *ordering customers*
G : **AL** *initially listed articles*
H : **WM** *initially available goods*
R : $P(\mathbf{SP})$ *available freight forwarders*

function symbols

f : **WA** \rightarrow **AR**
f : **WL** \rightarrow **AM**
(**_**)' : **AP** \rightarrow **AM**
(**_**)' : **AL** \rightarrow **AM**

variables

k: **KN**
x: **AR**
X, Y: **AL**
Z: $M(\mathbf{WA})$
t: **TE**
w: **WA**
s: **SP**
m, n, p: \mathbb{N}

properties

$(a, n)' = n[a]$ für $(a, n) \in \mathbf{AP}$
 $[p_1, \dots, p_n]' = p_1' + \dots + p_n'$ für
 $[p_1, \dots, p_n] \in \mathbf{AL}$

$\mathbf{G}' = \mathbf{f}(\mathbf{H})$

Terms of a signature Σ

$f(x)$

$f(a, g(y))$

Infinitely many terms

In a Σ -structure,
each term has a meaning!

a signature, Σ

ground sorts

KN *customers*

AR *articles*

WA *goods*

TE *dates*

SP *freight forwarders*

derived sorts

AP = **AR** $\times \mathbb{N}$ *article items*

AL = $M(\mathbf{AP})$ *article lists*

AM = $M(\mathbf{AR})$ *sets of articles*

WM = $M(\mathbf{WA})$ *sets of goods*

constant symbols

p1, p2: **AP** *ordered article positions*

B: $P(\mathbf{KN})$ *ordering customers*

G : **AL** *initially listed articles*

H : **WL** *initially available goods*

R : $P(\mathbf{SP})$ *available freight forwarders*

function symbols

f : **WA** \rightarrow **AR**

f : **WL** \rightarrow **AM**

(_'): **AP** \rightarrow **AM**

(_'): **AL** \rightarrow **AM**

variables

k: **KN**

x: **AR**

X, Y: **AL**

Z: $M(\mathbf{WA})$

t: **TE**

w: **WA**

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m, n, p: \mathbb{N}

properties

$(a, n)' = n[a]$ für $(a, n) \in \mathbf{AP}$

$[p_1, \dots, p_n]' = p_1' + \dots + p_n'$ für

$[p_1, \dots, p_n] \in \mathbf{AL}$

$G' = f(H)$

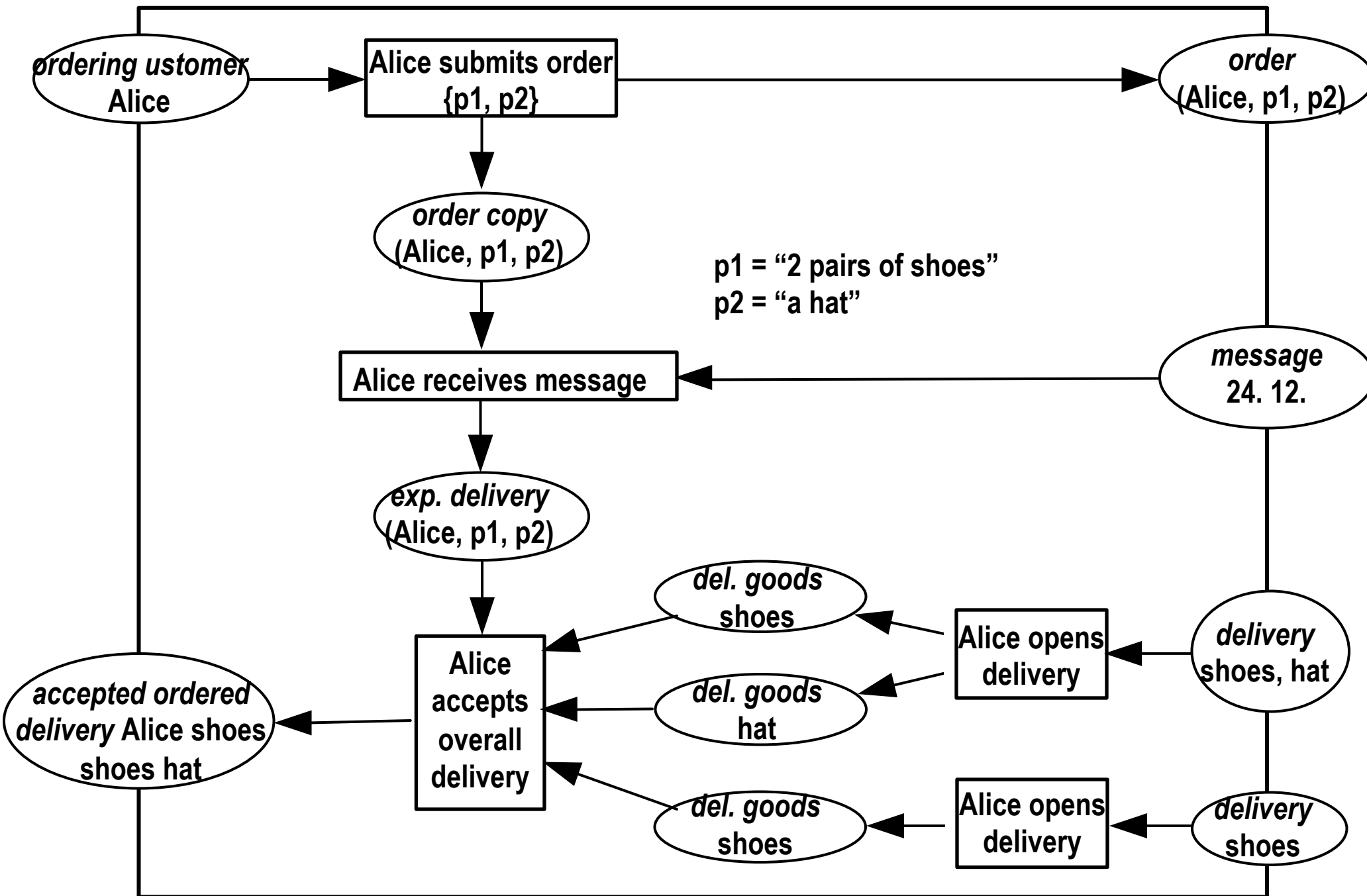
3. Dynamics

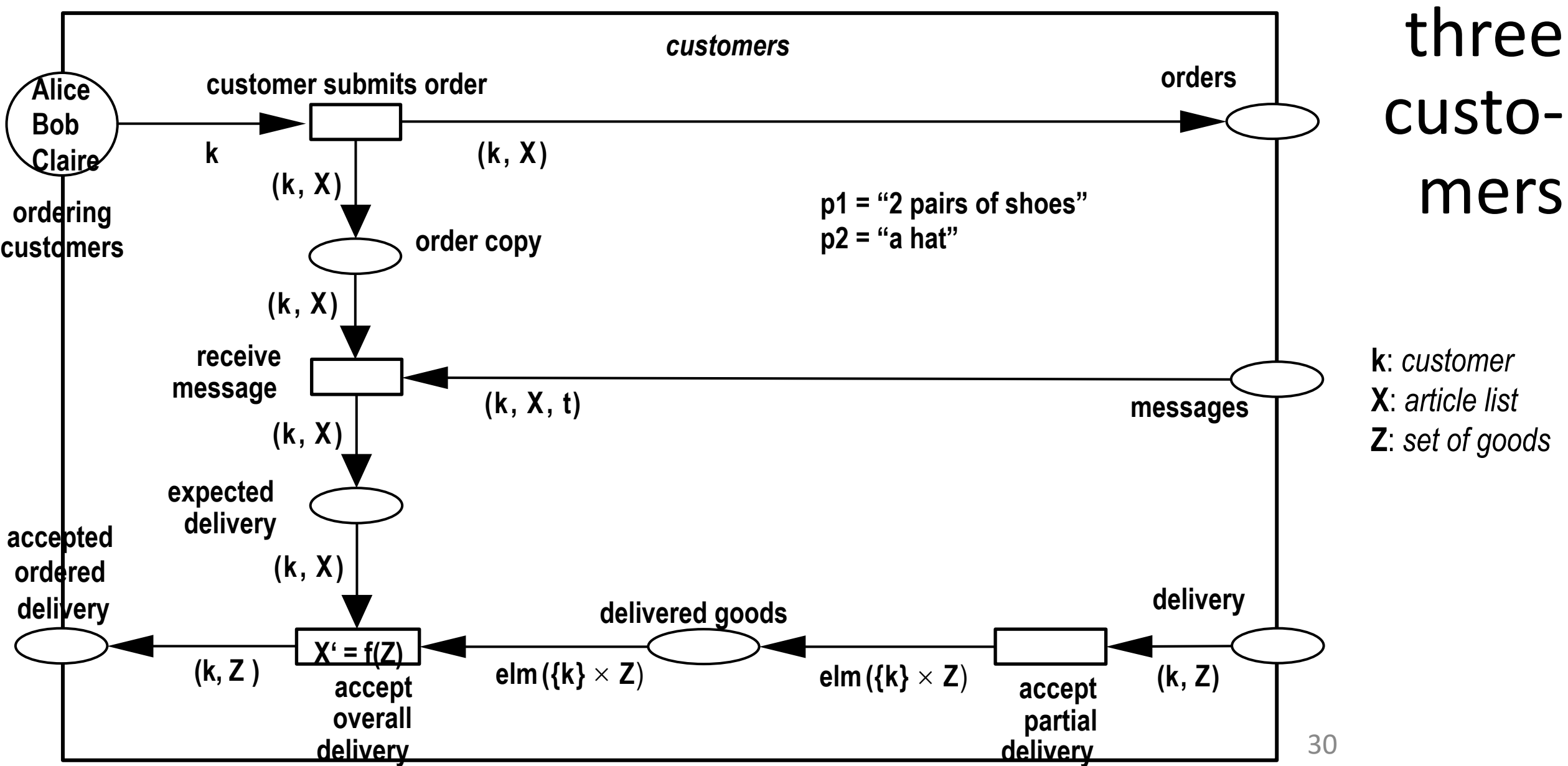
local states

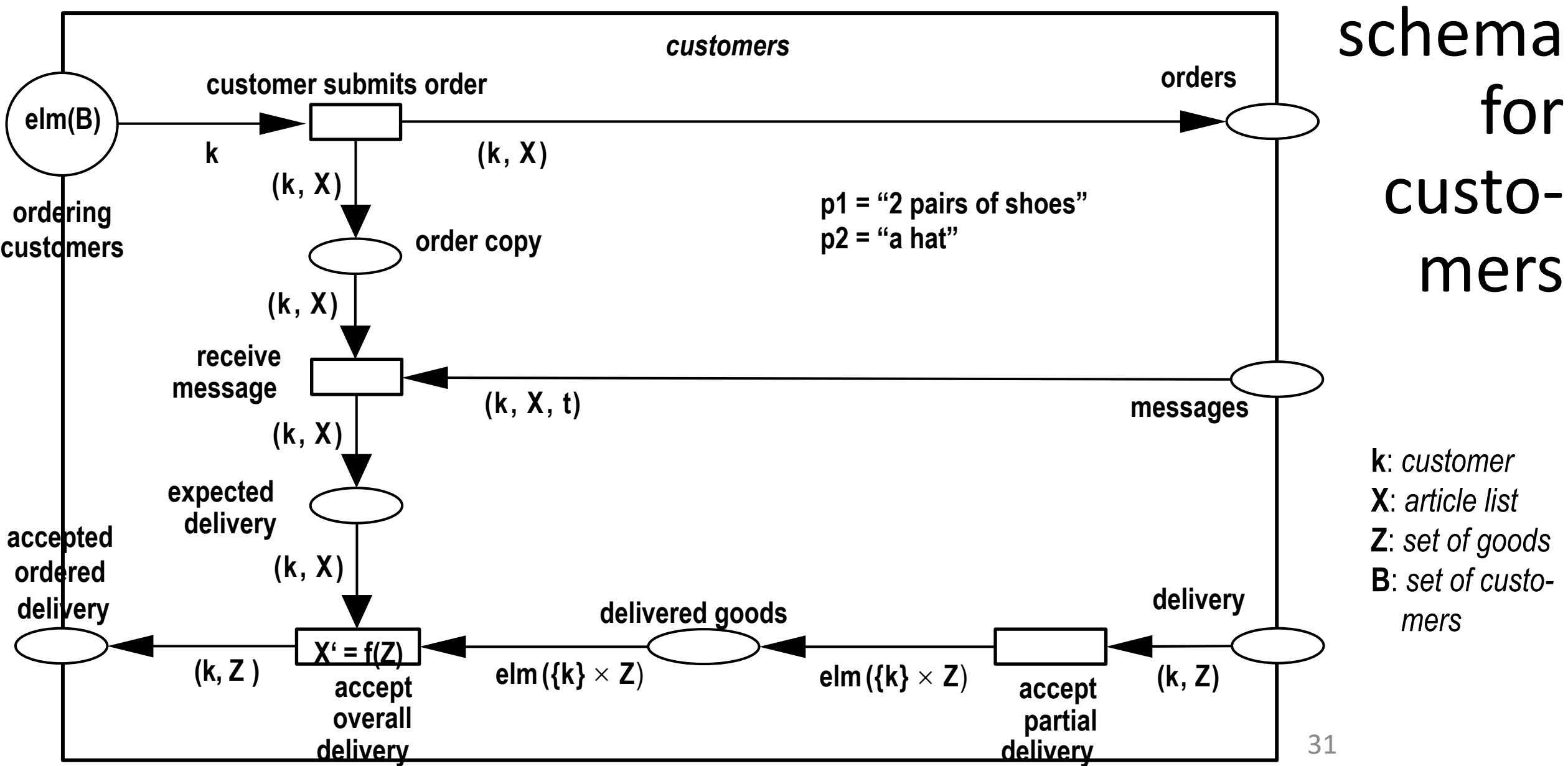
an event has local cause and effect

run of customer Alice

p1 = "2 pairs of shoes"
p2 = "a hat"

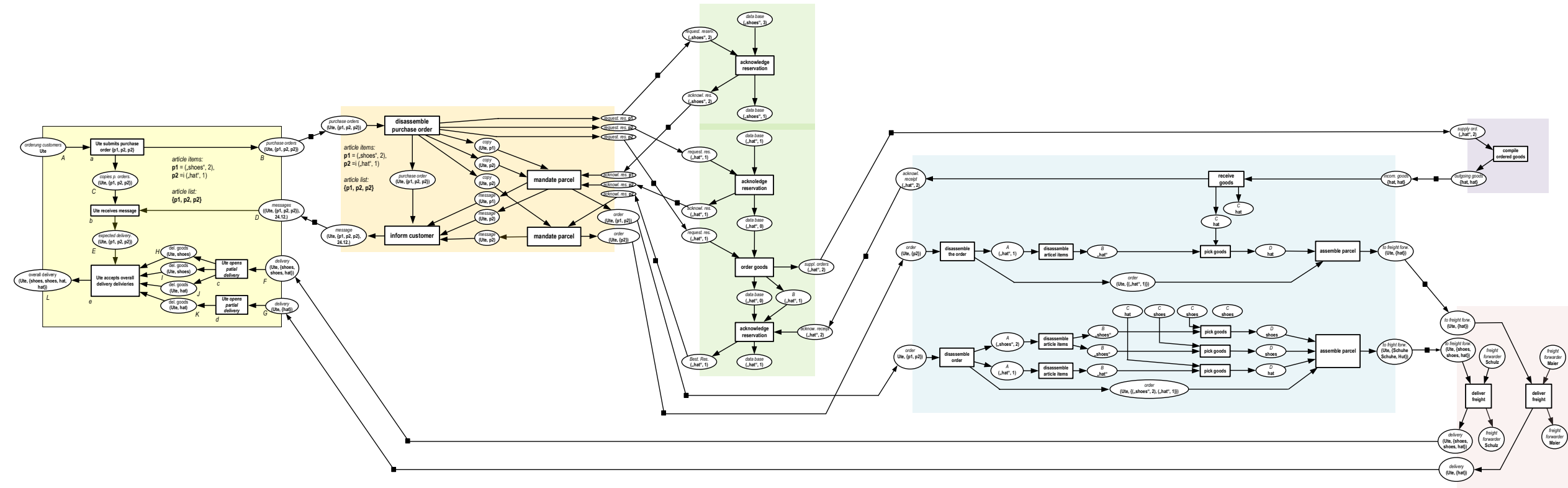






4. Put it all together

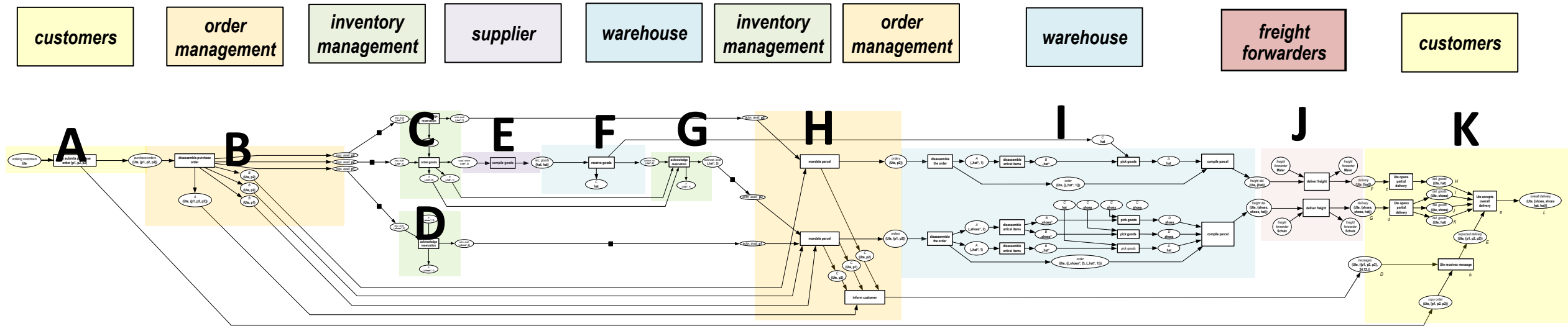
overall run of Alice's order



just write

customers • order management • inventory management • warehouse • supplier • freight forwarders

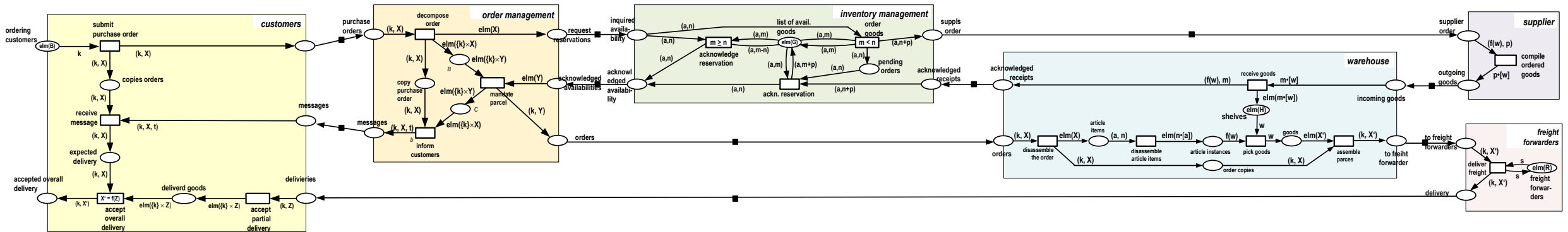
overall run of Alice's order represented from left to right



just write

A • B • C • D • E • F • G • H • I • J • K

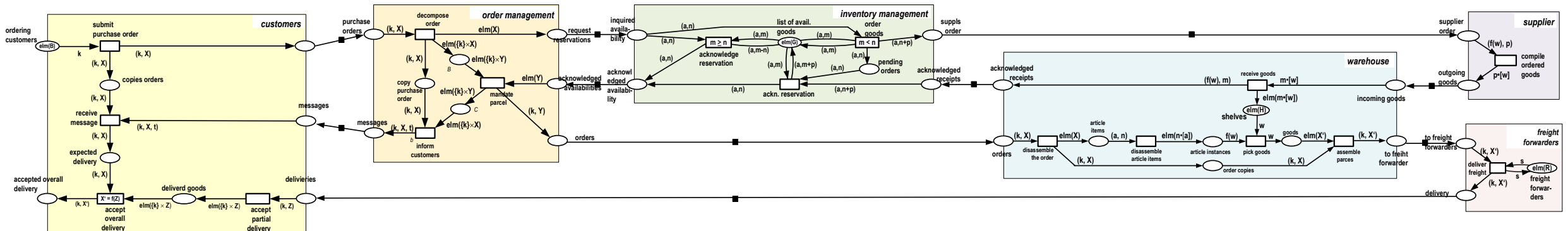
schema for interent shopping



just write

customers • order management • inventory management • warehouse • supplier • freight forwarders

schema for interent shopping



This is a formal model!

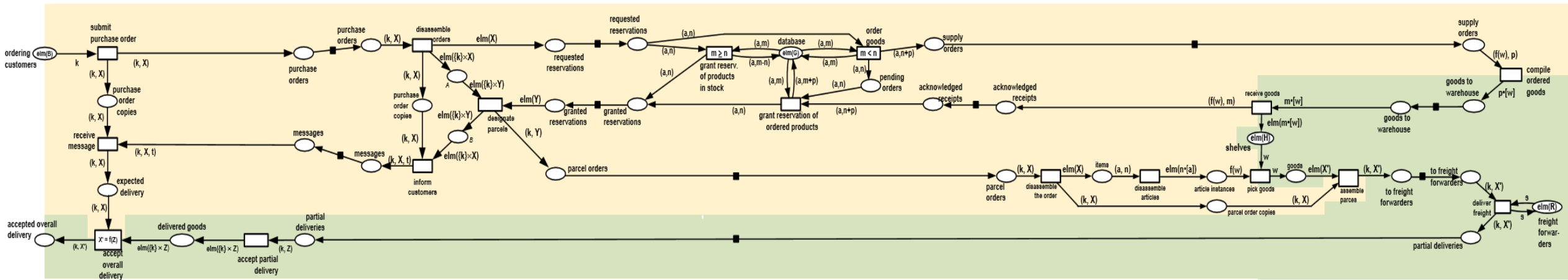
You can formally prove:

- Each order is eventually served.
- A good is delivered only if it has been ordered before.
- Each requested reservation is eventually acknowledged.
- Each mandated article has a good in the shelves.



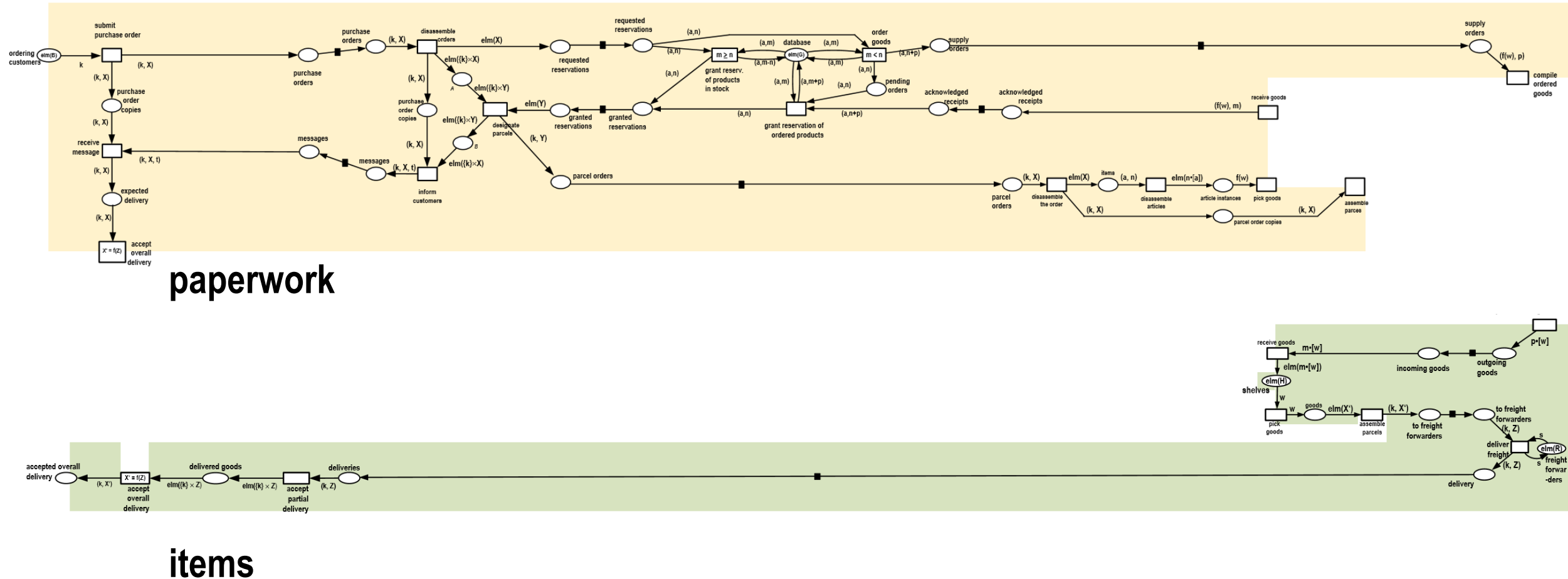
alternative refinement

business = paperwork • items



alternative refinement

business = paperwork • items



How to model the internet of everything

classical informatics

modules and composition:
merge „equal“ interface
elements

statics (data, items):
symbolic representation

dynamics: steps

- classical computer science
- jumps in the right direction
 - but falls short

... yes, but ...

yes, however not
one interface

yes, however not
with symbol *chains*

yes, however, not
global states and steps



... adjusted

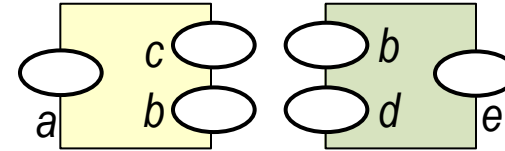
but two!

but with terms
over a signature!

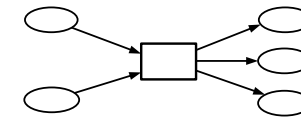
but local ones!

adjusts this!

... such as



$f(x, g(a, y))$



... technically

composition
calculus

predicate logic,
algebraic spec.

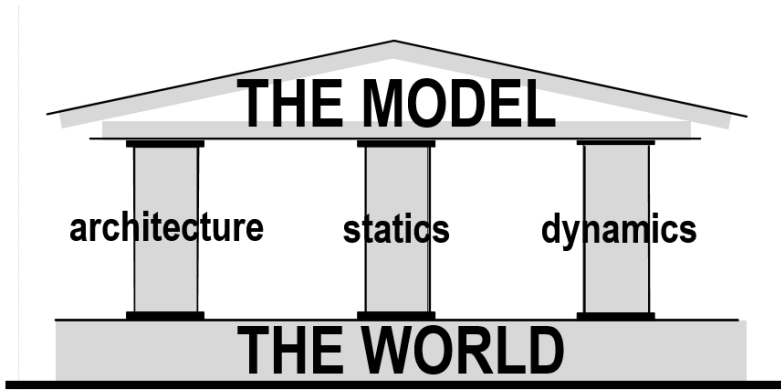
Petri nets

Don't hesitate!

Join the  research program and
development project

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do it with



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