## Next generation modeling in the era of cyber-physical systems

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## First generation modeling anno 1993: Hasso Plattner and Klaus Besier pose with the SAP ERP Reference Model

Peter



#### **DigitalOps Combines Three Domains**



Note: Organizations are moving to combine these domains, rather than dealing with them separately.



Part I examples

- 1. architecture
- 2. single behavior
- 3. elementary systems
- 4. items and data

Part II A glimpse at concepts:

- The three HERAKLIT pillars
- 5. architecture: Two-faced modules
- 6. dynamics: steps: from requirements to models

7. statics: Breathing live into logic: structures and signatures

Part III A big case study: an apetizer

## Part I Examples 1. Architecture

### The four activities of a bakery

#### Wolfgang

- A big system consists of *modules*.
- A module has an *interface*.

- An interface consists of *gates*.
- Each gate is *labeled*.





# Part I Examples

The four activities of a bakery

bake • supply to aide





## Part I Examples <sup>T</sup> 1. Architecture

The four activities of a bakery

bake • supply to aide • move to shop





# Part I Examples The of a 1. Architecture

The four activities of a bakery

bake • supply to aide • move to shop





# Part I ExamplesThe four activities<br/>of a bakery1. Architecture

bake • supply to aide • move to shop • sell





# Part I ExamplesThe four activities<br/>of a bakery1. Architecture

bake • supply to aide • move to shop • sell



# Part I ExamplesThe three staff<br/>of a bakery1. Architecture





## Part I Examples 2 behavior

Peter



Remember the modules









### Example: modules of staff



#### move to shop vendor shelf empty shelf empty

#### Remember the modules

#### **Operational behavior:**







#### Example: modules of staff





#### Example: modules of staff









#### Concurrency is not transitive!





Wolfgang







![](_page_24_Figure_1.jpeg)

![](_page_24_Figure_2.jpeg)

![](_page_24_Figure_3.jpeg)

![](_page_24_Figure_4.jpeg)

![](_page_25_Figure_1.jpeg)

baker system • aide system 1 •
vendor system • vendor system 1

versus

baker system • aide system 1 •
vendor system 1 • vendor system

![](_page_25_Figure_5.jpeg)

![](_page_26_Figure_1.jpeg)

baker system • aide system 1 •
vendor system • vendor system 1

versus

baker system • aide system 1 •
vendor system 1 • vendor system

![](_page_26_Figure_5.jpeg)

![](_page_27_Figure_1.jpeg)

#### Peter

## 4. items and data

![](_page_28_Figure_2.jpeg)

recent supply: descriptions,

*variables* x: pastries y: descriptions

### 4. items and data

![](_page_29_Figure_1.jpeg)

### Pause

## Part II a glimpse at concepts

#### Wolfgang

![](_page_31_Figure_2.jpeg)

Part II A glimpse at concepts:

The three HERAKLIT pillars

5. architecture: Two-faced modules

6. dynamics: steps: from requirements to models

7. statics: Breathing live into logic: structures, signatures and schamata

Part III A big case study: an apetizer

#### **Theoretical informatics**

Given an alphabet  $\Lambda \{\alpha, \beta, \gamma\}$ .

Canonical constructs:

- word over  $\Lambda = \beta \gamma \beta \alpha \alpha$
- Set of all words over  $\Lambda$ , written  $\Lambda^*$
- Composition of words:  $\alpha\beta\beta \bullet \beta\alpha\gamma \bullet \beta\gamma\gamma$ =  $\alpha\beta\beta\beta\alpha\gamma\beta\gamma\gamma$

Monoid ( $\Lambda^*$ , •,  $\epsilon$ ):

THE formal fundament of computing.

#### HERAKLIT

Given an alphabet  $\Lambda = \{\alpha, \beta, \gamma\}.$ 

Canonical constructs:

- Modules with gate labels in  $\Lambda$
- Set of all modules over  $\Lambda$ , written  $\Lambda M$
- Composition of modules M N

Monoid ( $\Lambda M$ , •,  $\epsilon$ ): THE formal fundament of modeling.

![](_page_33_Figure_1.jpeg)

A module is

- a graph

![](_page_34_Figure_1.jpeg)

A module is

- a graph
- With two distinguished sets of nodes (left and right *interface*) *"gates"*

![](_page_35_Figure_1.jpeg)

A module is

- a graph,
- With two distinguished sets of nodes (left and right *interface*) "gates".
- Each gate is labeled.

![](_page_36_Figure_1.jpeg)

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Here a second module,  $N_0$ .

A module is

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Composition of  $M_0$  and  $N_0$ 

![](_page_37_Figure_7.jpeg)

![](_page_37_Figure_8.jpeg)

- Any two modules can be composed, resulting in a module.
- Compsosition is associative:
- $\mathsf{L} \bullet (\mathsf{M} \bullet \mathsf{N}) = (\mathsf{L} \bullet \mathsf{M}) \bullet \mathsf{N}$
- For the empty module  $\epsilon$  holds:

 $\mathsf{M} \bullet \epsilon = \epsilon \bullet \mathsf{M} = \mathsf{M}$ 

- A gate may lie in the left as well as in the right interface

![](_page_38_Figure_7.jpeg)

![](_page_38_Figure_8.jpeg)

![](_page_39_Picture_1.jpeg)

Peter

![](_page_40_Picture_1.jpeg)

In the case of the fan off, when you turn on the light, after some time, the fan will start running. In this situation, if you turn off the light, the fan continues running for some time. Hence, in the case of the fan off, when you turn on and off the light quickly, the fan will not start running at all. And in the case of the fan on, when you turn off and on the light quickly, the fan will continuously run.

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![](_page_42_Figure_2.jpeg)

![](_page_43_Picture_0.jpeg)

![](_page_43_Picture_1.jpeg)

![](_page_44_Figure_0.jpeg)

# 7. Breathing live into logic Varying propositions

#### Wolfgang

The notion of *proposition* 

Aristotle:Petri:Always truesometimes true

 $e = mc^2$  A bread lies on the shelve

This is *not* temporal logic! TL: a proposition abstracts a global state Petri: a proposition *IS* a (locally confined) state

# 7. Breathing live into logic example

![](_page_46_Figure_1.jpeg)

## 7. Breathing live into logic Two propositions

![](_page_47_Figure_1.jpeg)

# 7. Breathing live into logic three propositions

![](_page_48_Figure_1.jpeg)

# 7. Breathing live into logic Predicates and parameterized events

![](_page_49_Figure_1.jpeg)

Variable x: pastry

![](_page_50_Figure_0.jpeg)

## 7. Structures, signatures, and schemata

![](_page_51_Figure_1.jpeg)

#### Remember ... Three pastries

<i>domains</i> pastries = {bread, cake, pie} descriptions = {"bread", "cake", "pie"}	<i>function</i> next: descriptions → pastries next("bread") = cake next("cake") = pie next("pie") = bread	

*constants* "pie": descriptions predicates on counter, aide with pastry, on shelf: pastries recent supply: descriptions, bakery with three pastries

#### function

des: pastries → descriptions des(bread) = "bread" des(cake) = "cake" des(pie) = "pie"

*propositions* aide free, shelf empty

### 7. Structures, signatures, and schemata

![](_page_52_Figure_1.jpeg)

![](_page_52_Figure_2.jpeg)

five pastries

*domains* pastries = {bread, cake, pie, rol, biscuit} descriptions = {,,bread", ,,cake", ,,pie", ,,rol", ,,biscuit"}

#### function

des: pastries → descriptions des(bread) = "bread" des(cake) = "cake" des(pie = "pie" des(rol) = "rol" des(biscuit) = "biscuit"

#### *constants* "bread": descriptions

predicates

on counter, aide with pastry, on shelf: pastries recent supply: descriptions

#### function

next: descriptions → pastries next("bread") = cake next("cake") = pie next("pie") = rol next("rol") = biscuit next("biscuit") = bread

propositions aide free, shelf empty

## 7. Structures, signatures, and schemata

![](_page_53_Figure_1.jpeg)

Schema: Any set of pastries

		<i>signature</i> bakery
<u>domains</u> pastries descriptions	<u>function</u> next: descriptions — pastries	<u>function</u> des: pastries → descriptions
<u>constants</u> <u>p: descriptions</u>	<u>predicates</u> on counter, aide with pastry, on shelf: p recent supply: descriptions	propositions astries aide free, shelf empty

### Part III A big case study: an apetizer

abstract view

![](_page_55_Figure_0.jpeg)

## Part III A big case study: an apetizer abstract view

![](_page_56_Figure_0.jpeg)

#### customers • business • freight forwarders • suppliers

![](_page_57_Figure_0.jpeg)

customers • business • freight forwarders • suppliers 59 where business = order management • inventory management • warehouse

![](_page_58_Figure_0.jpeg)

the business = customers • [retailer] • [supplier] • [freight forwarder]

## overall run of Alice's order

![](_page_59_Figure_1.jpeg)

#### just write <u>customers</u> • <u>order management</u> • <u>inventory management</u> • <u>warehouse</u> • <u>supplier</u> • <u>freight forwarders</u>

## How to represent composition of run snippets?

![](_page_60_Figure_1.jpeg)

just write

 $A \bullet B \bullet C \bullet D \bullet E \bullet F \bullet G \bullet H \bullet I \bullet J \bullet K$ 

## schema for interent shopping

![](_page_61_Figure_1.jpeg)

#### just write

customers • order management • inventory management • warehouse • supplier • freight forwarders

## alternative refinement

#### business = paperwork • items

![](_page_62_Figure_2.jpeg)

## alternative refinement

#### business = paperwork • items

![](_page_63_Figure_2.jpeg)

items

#### Summing up: central ideas of HERAKLIT

classical computer science	yes, but	adjusted	such as	technically
modules and composition: merge "equal" interface elements	yes, however not one interface	but two!		composition calculus
statics (data, items): symbolic representation	yes, however not with symbol <i>chains</i> ("strings")	but with terms over a signature!	f(x, g(a,y))	predicate logic, algebraic specification
dynamics: steps	yes, however, not global states and steps	but local ones!		Petri nets
<ul><li>classical computer science</li><li>jumps in the right direction</li><li>but falls short</li></ul>	HERAKLIT adju	sts this! Programm Modeling	ing is about sym is about the digit	bol crunching. al world!