

Trust by design

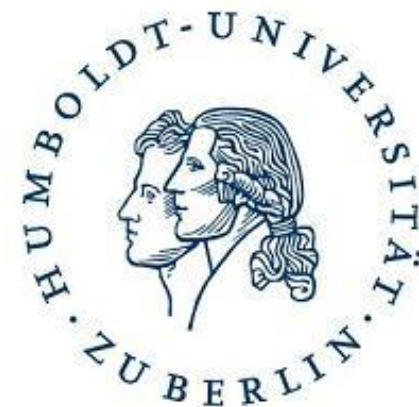
- *in praise of modularization:
a case study*

Wolfgang Reisig

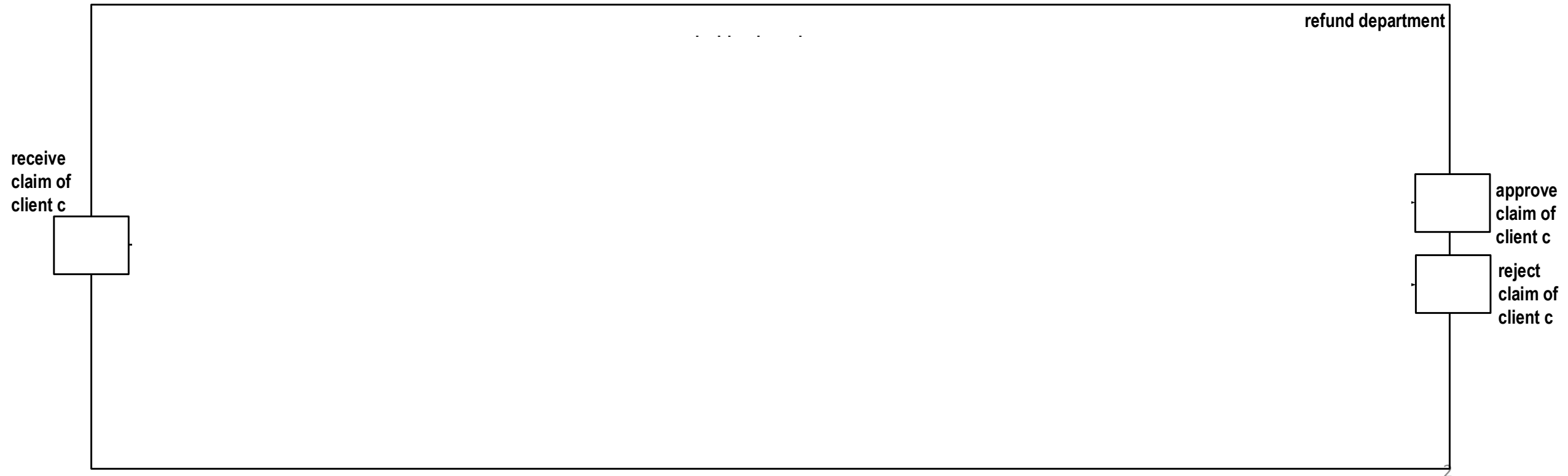
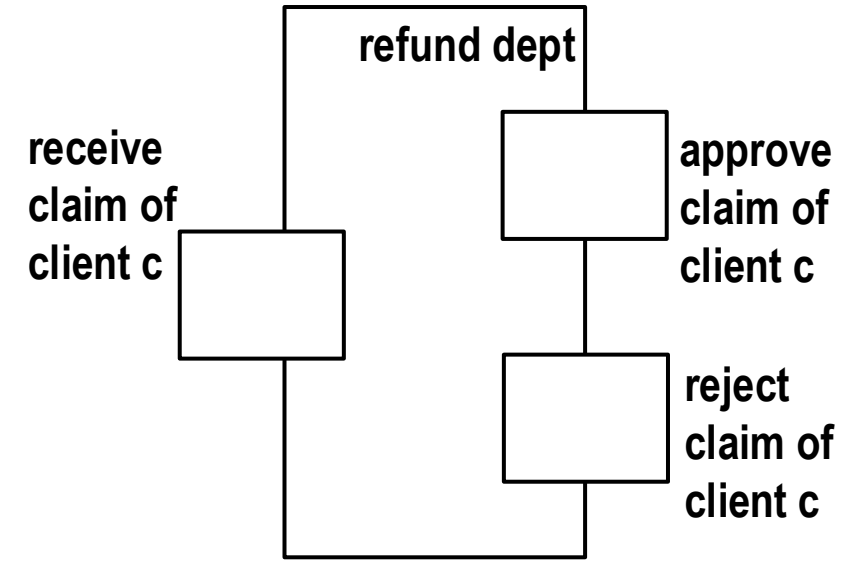
*Humboldt-Universität
zu Berlin
Germany*

Peter Fettke

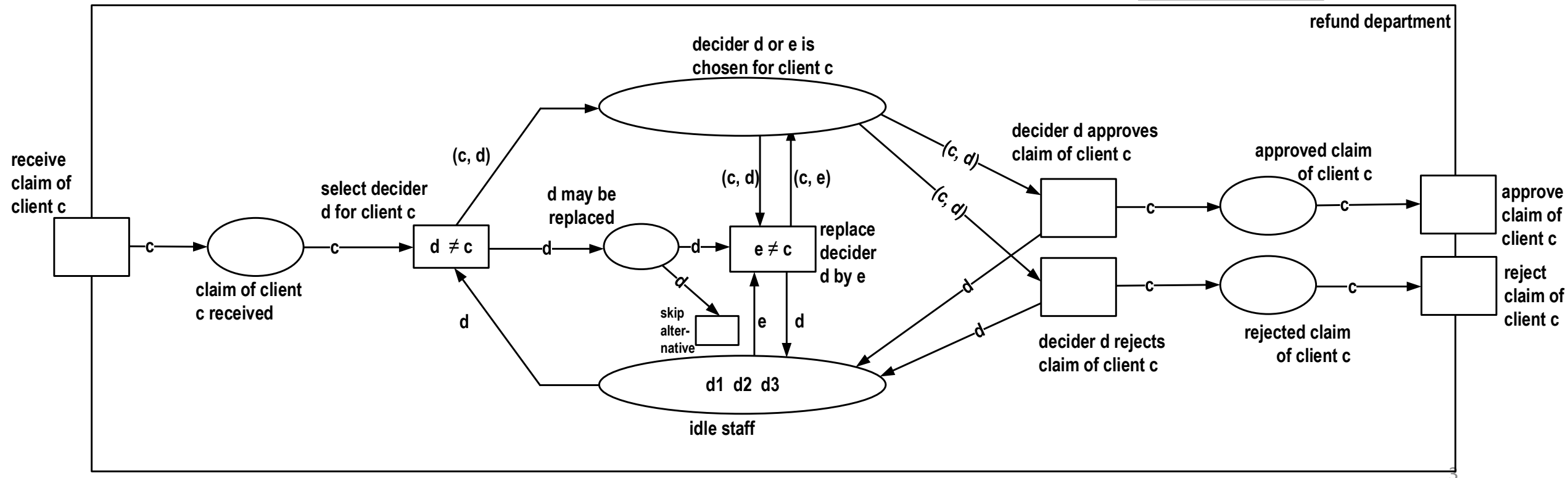
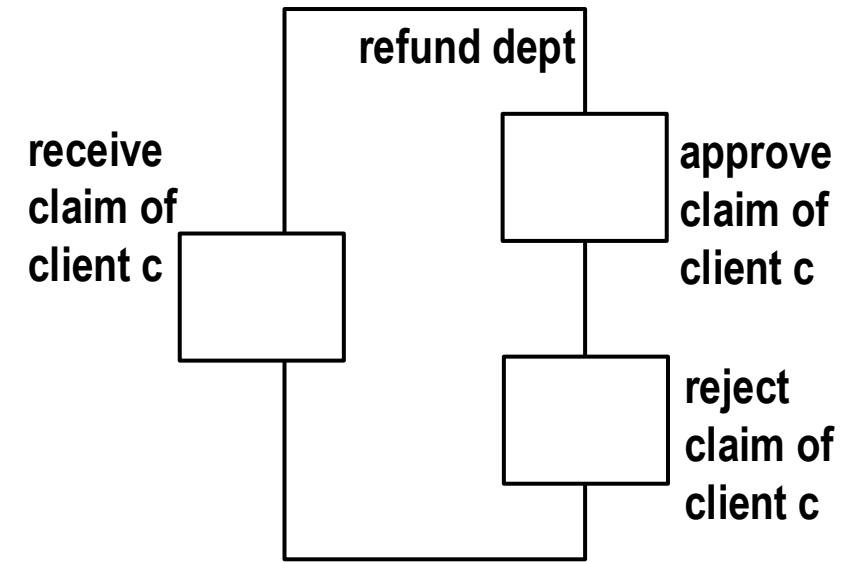
*German Research Center
for Artificial Intelligence
(DFKI) and Saarland
University, Saarbrücken,
Germany*



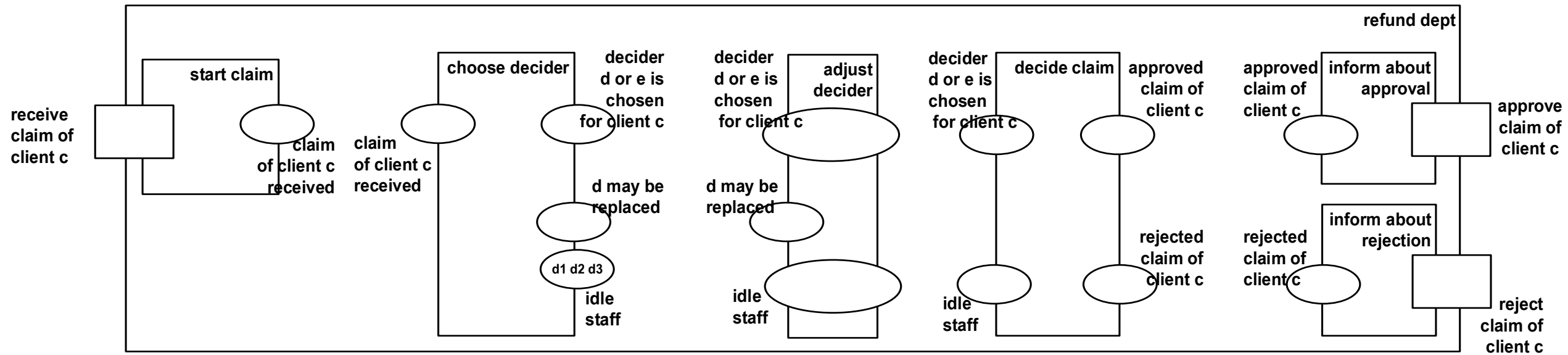
most abstract view



... grasp this in one go?



better: six modules of the system model

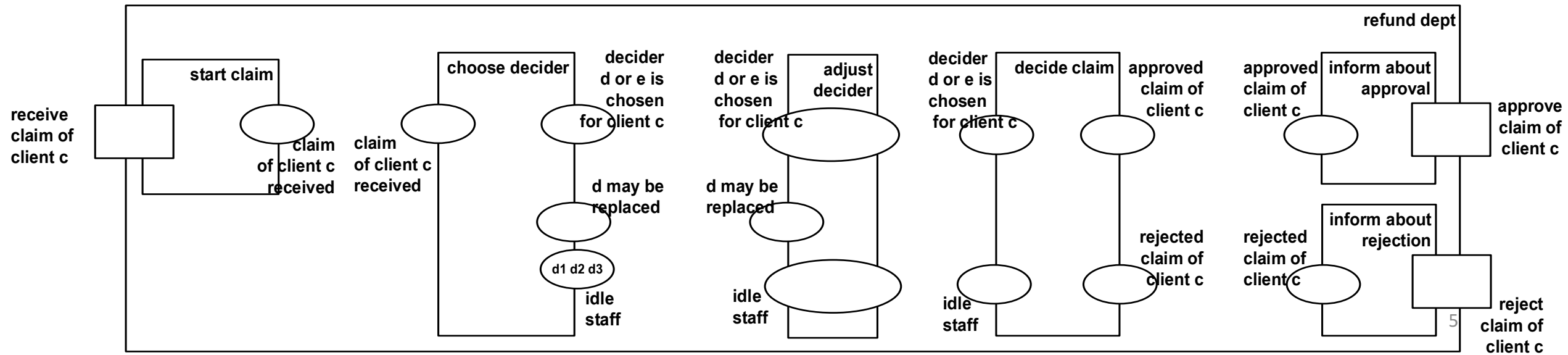


This talk

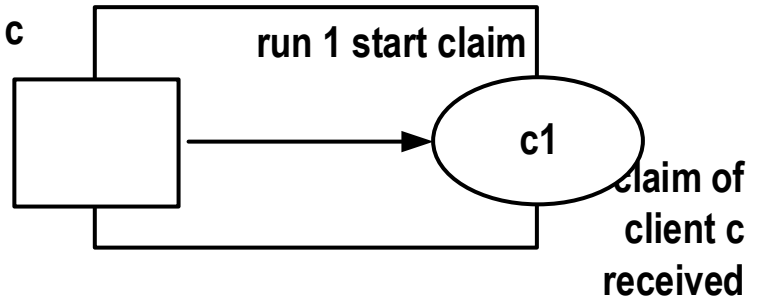
Part I Run modules and typical runs

Part II System modules

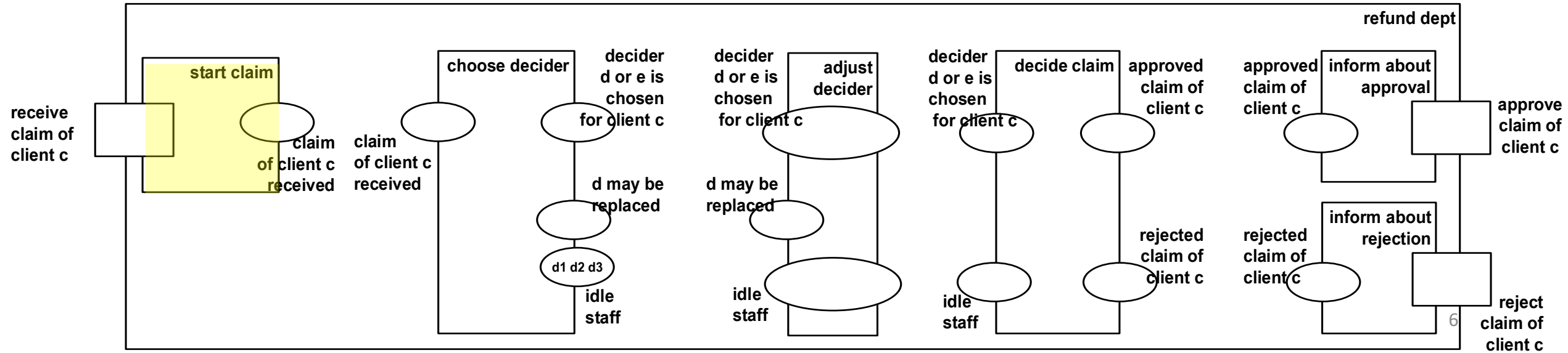
Part III Liveness



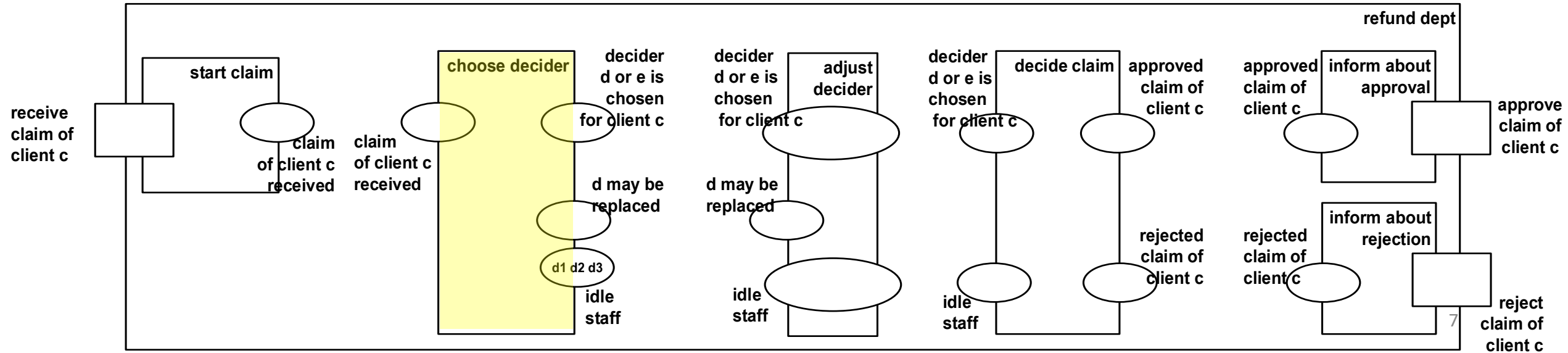
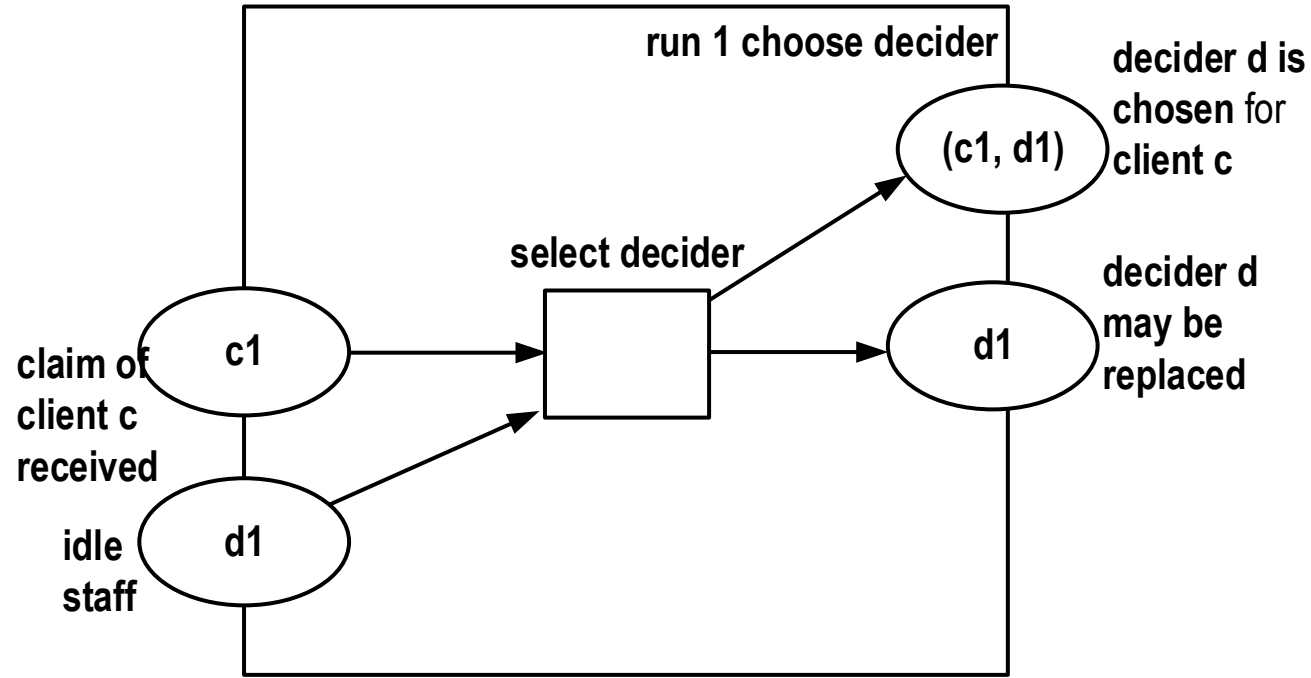
receive
claim of
client c



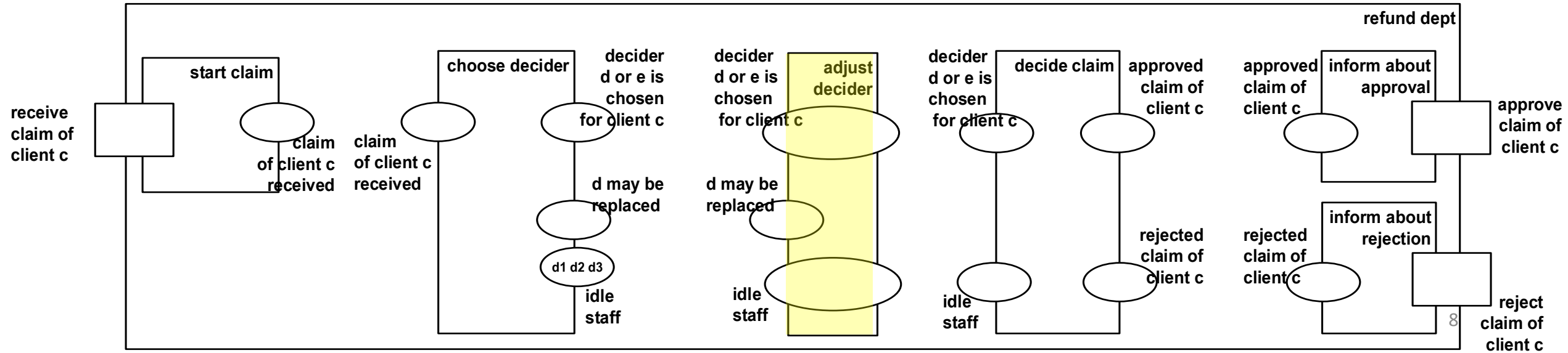
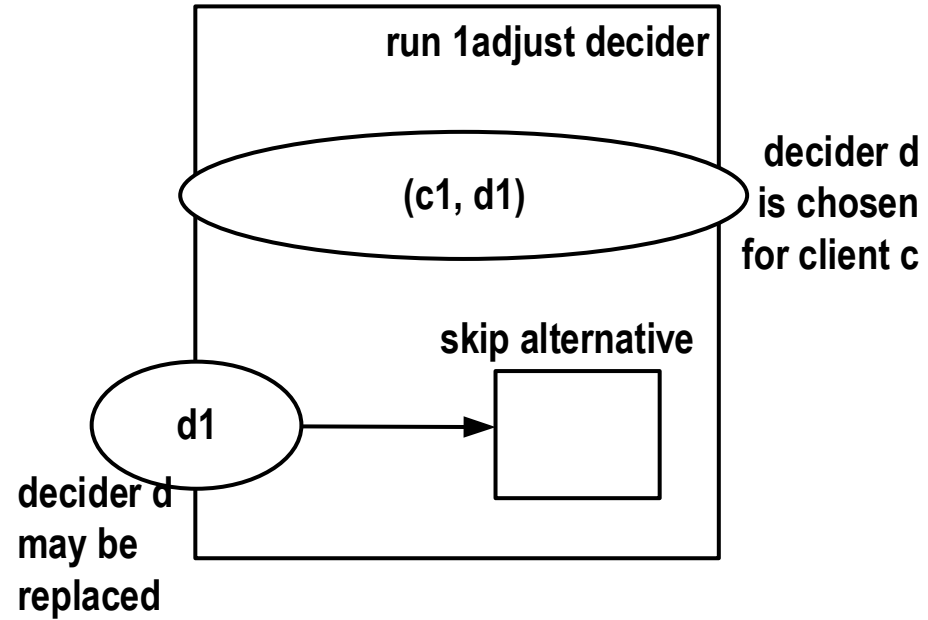
Part I run modules



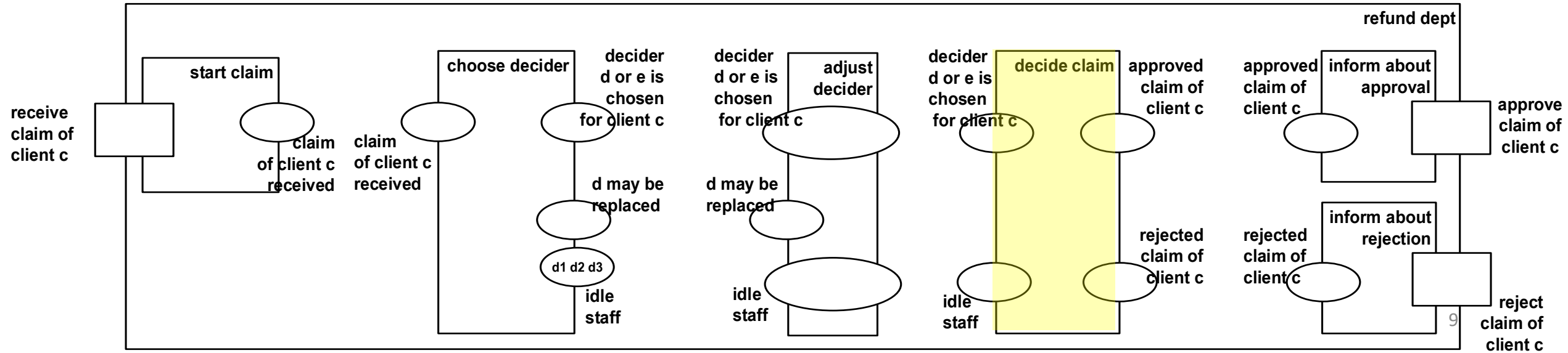
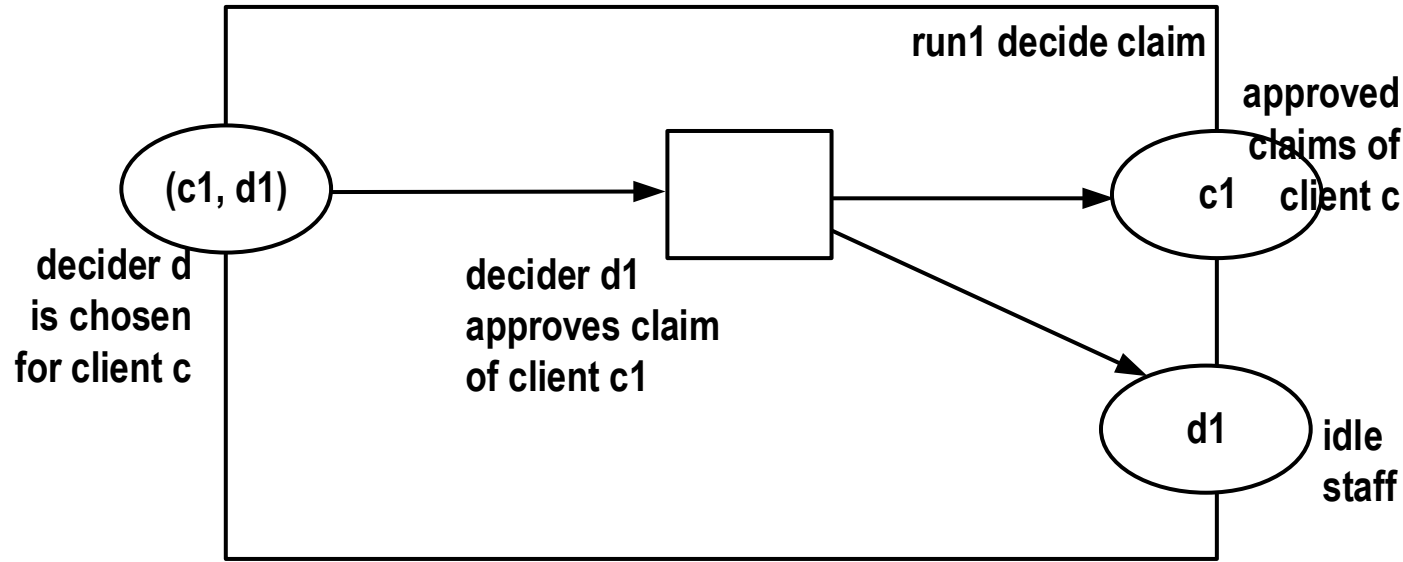
Part I run modules



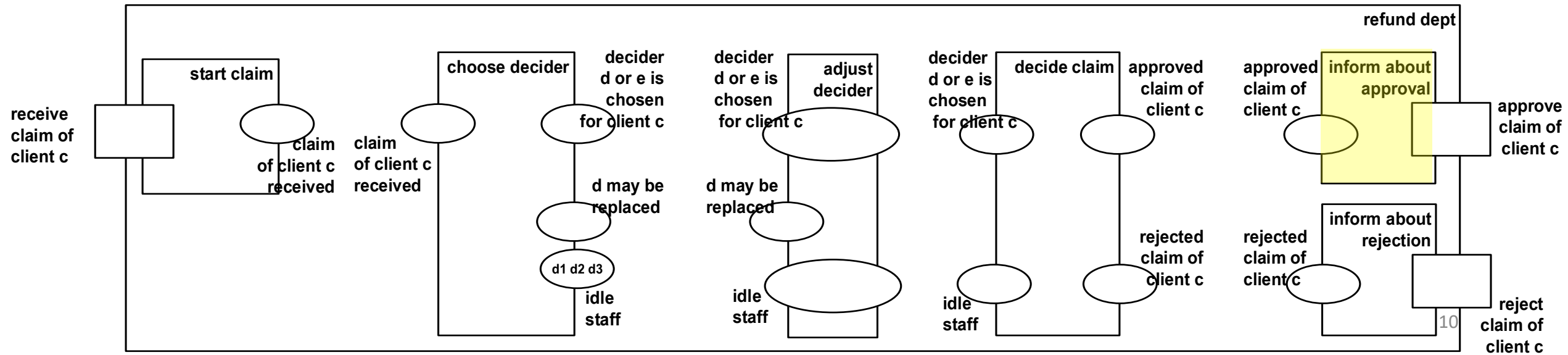
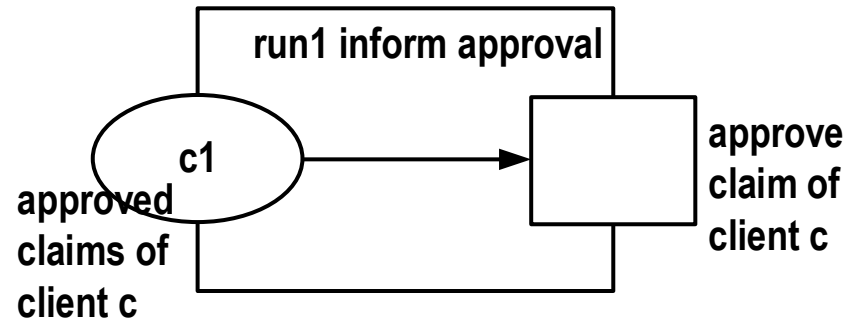
Part I run modules

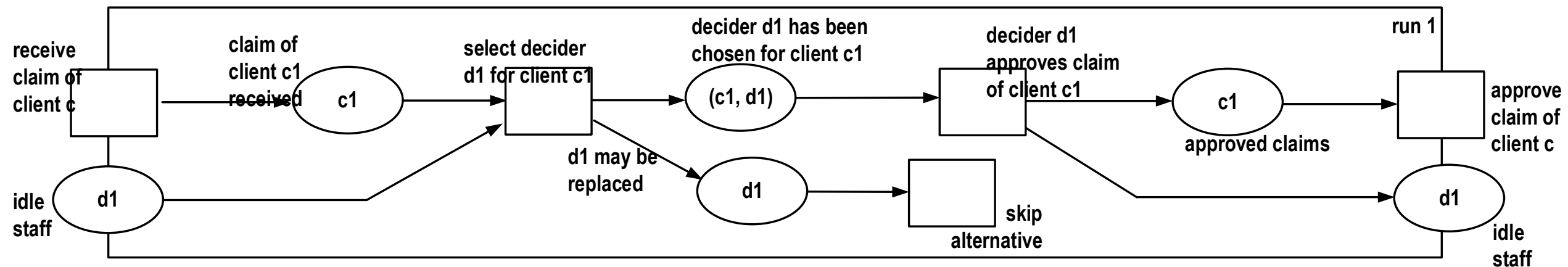


Part I run modules



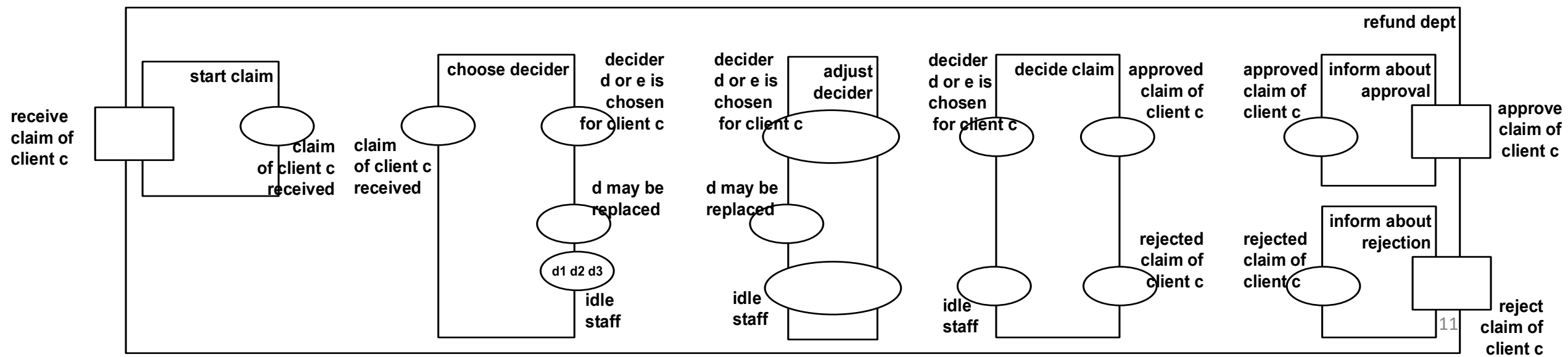
Part I run modules

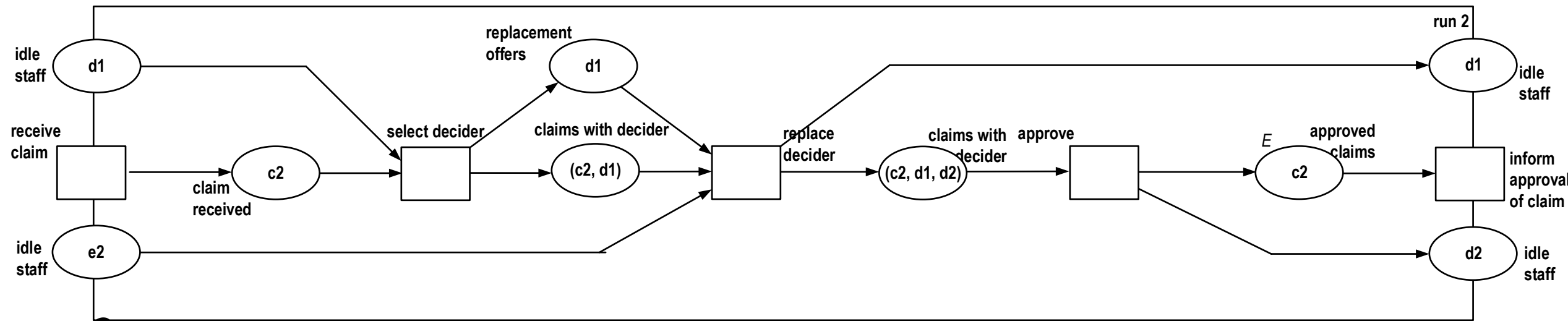




run1 =

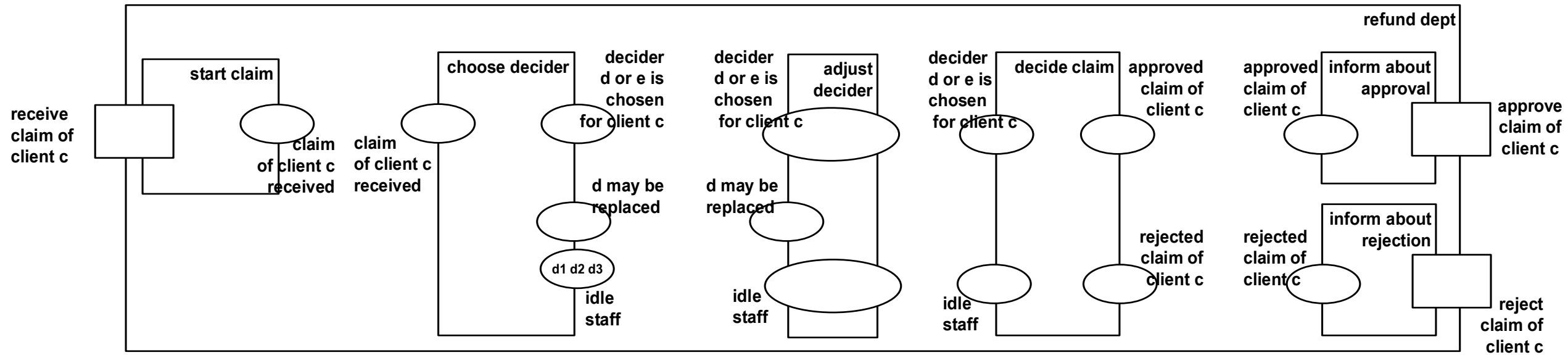
run 1 start claim • run1 choose decider • run1 adjust decider • run1 decide claim • run1 inform about approval



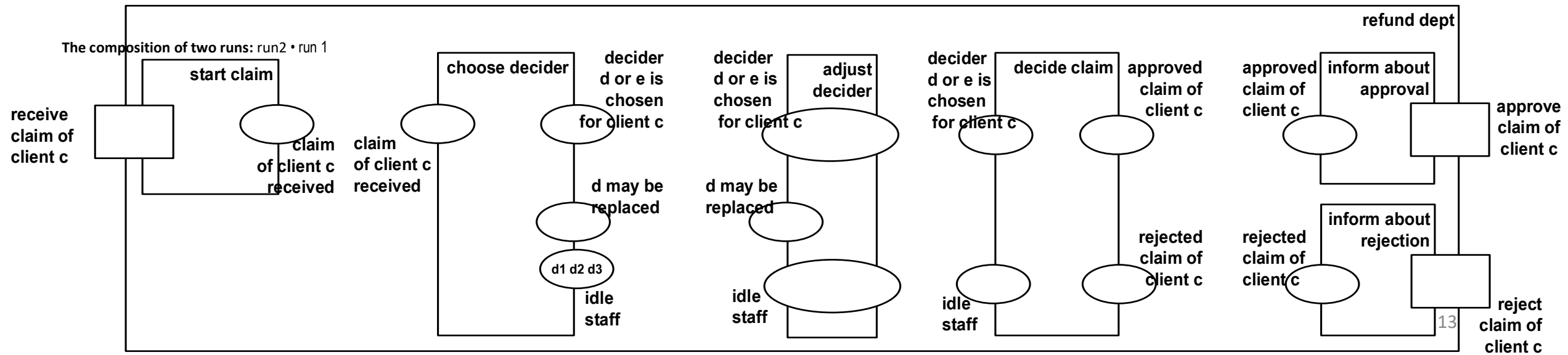
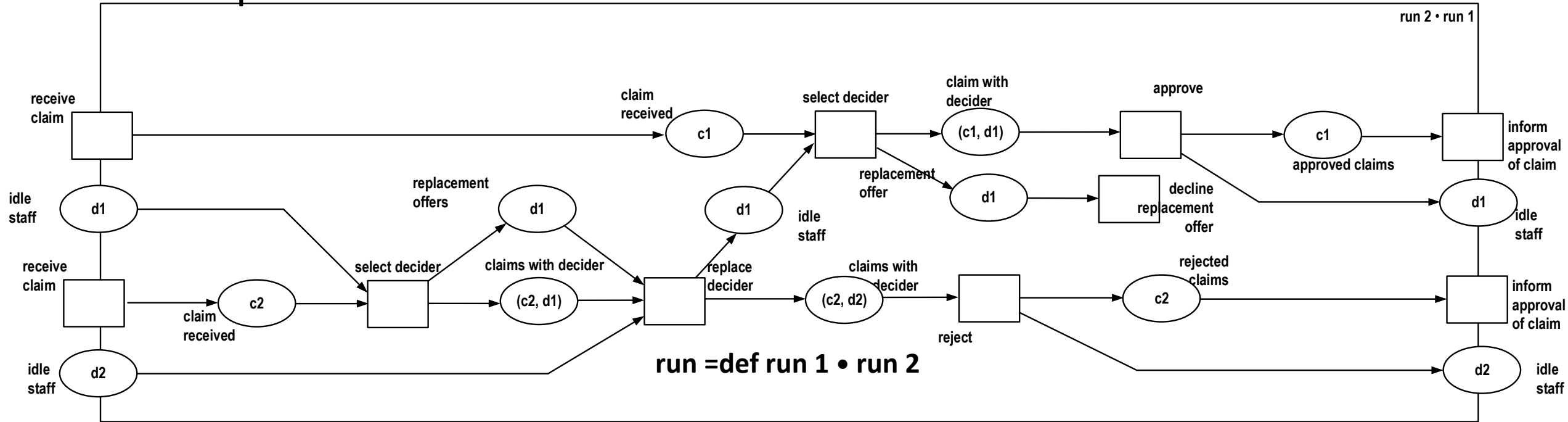


run 2 =

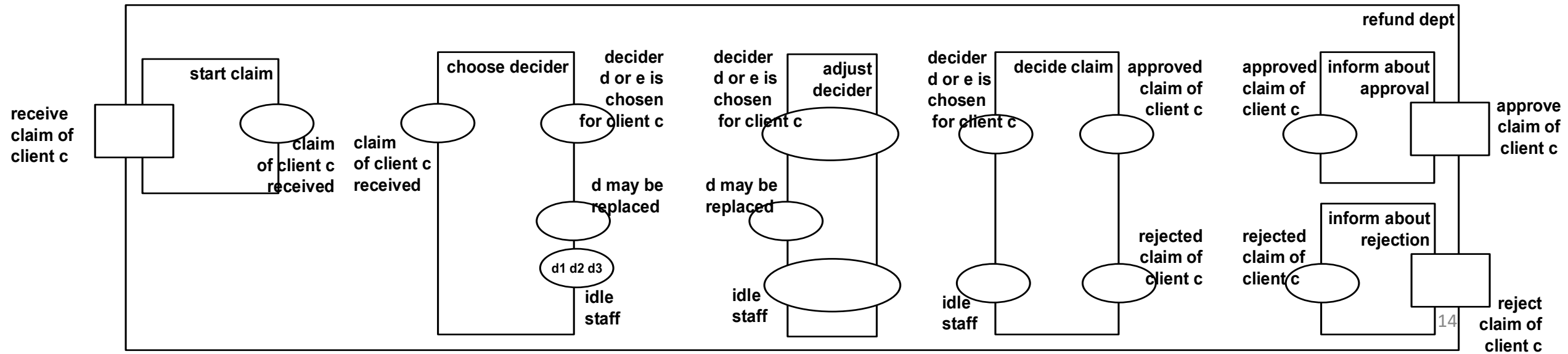
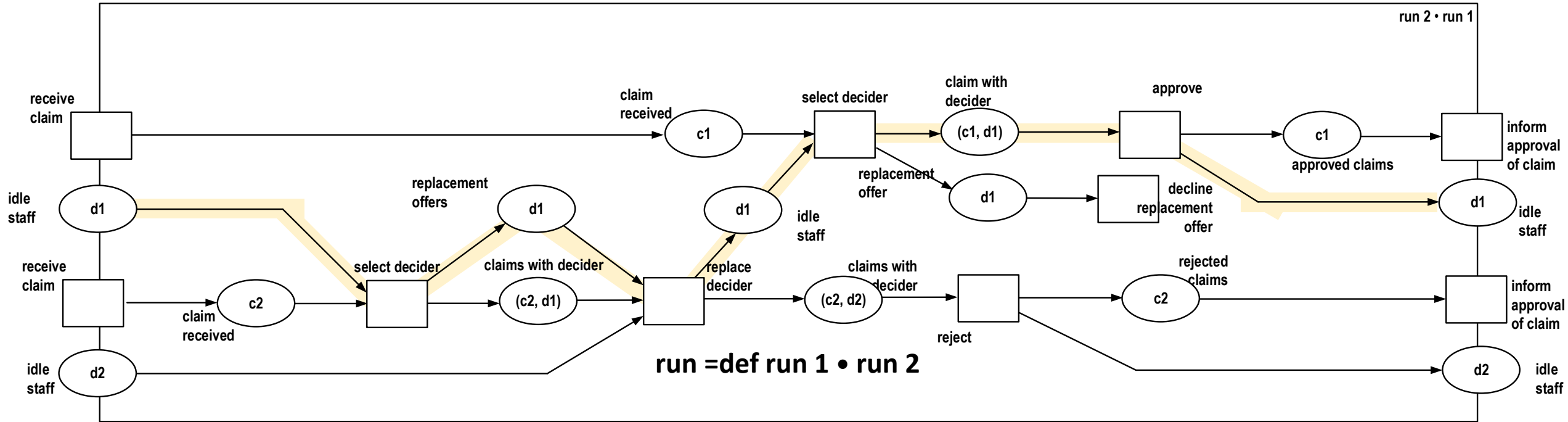
run 2 start claim • run2 choose decider • run2 adjust decider • run2 decide claim • run2 inform about rejection



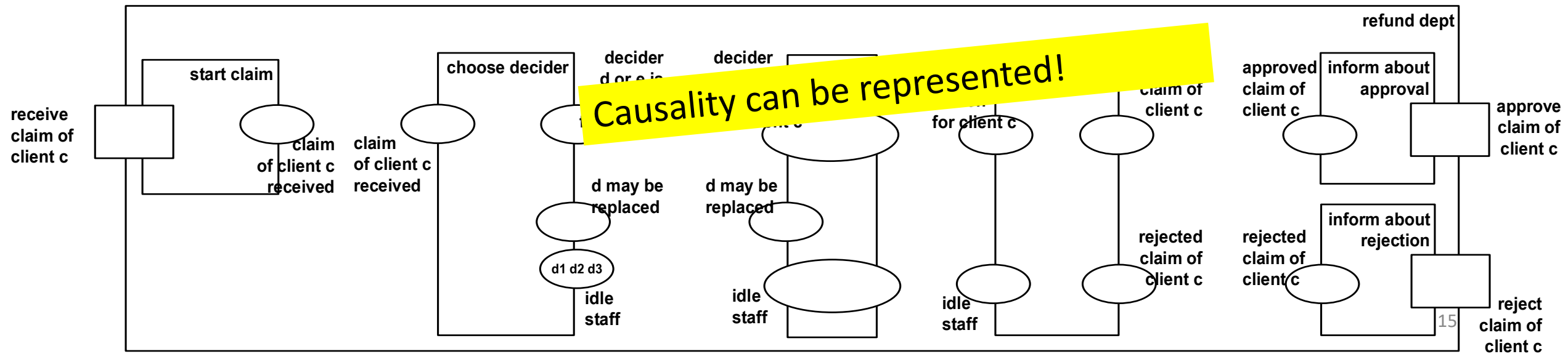
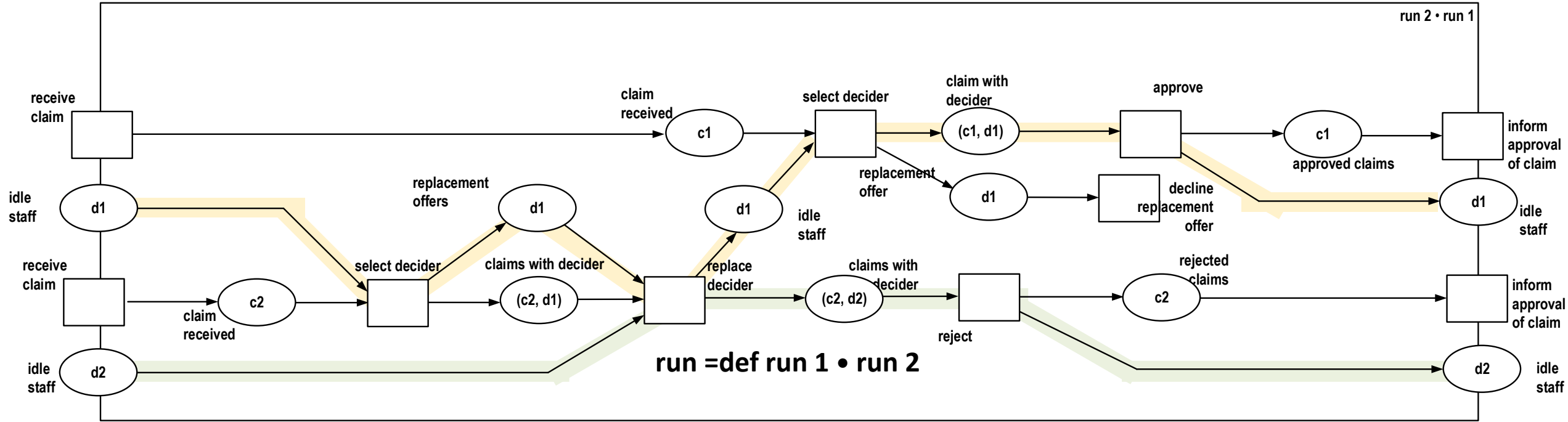
composition of both runs



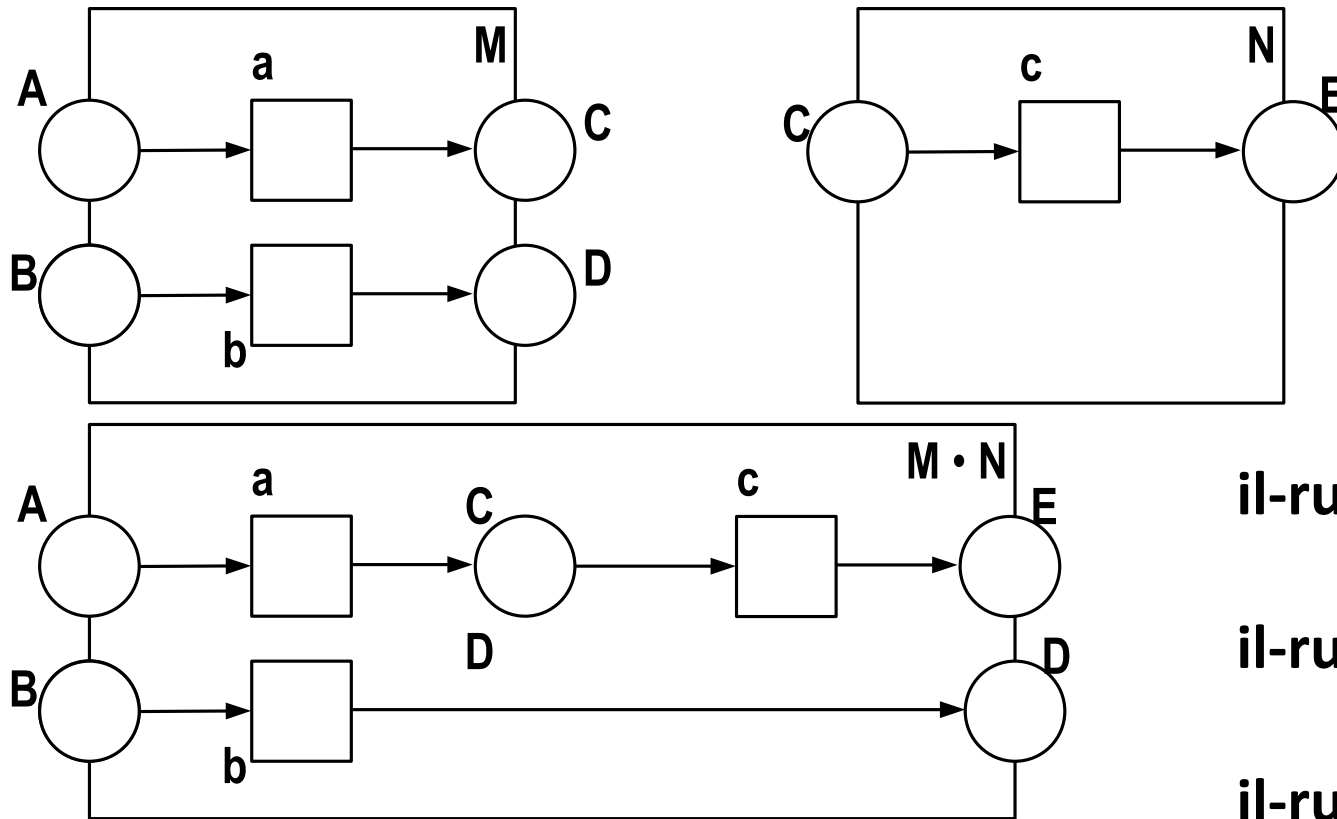
the lifeline of d1



the lifeline of d1 and of d2



runs are compositional!



interleaved runs are
not compositional!

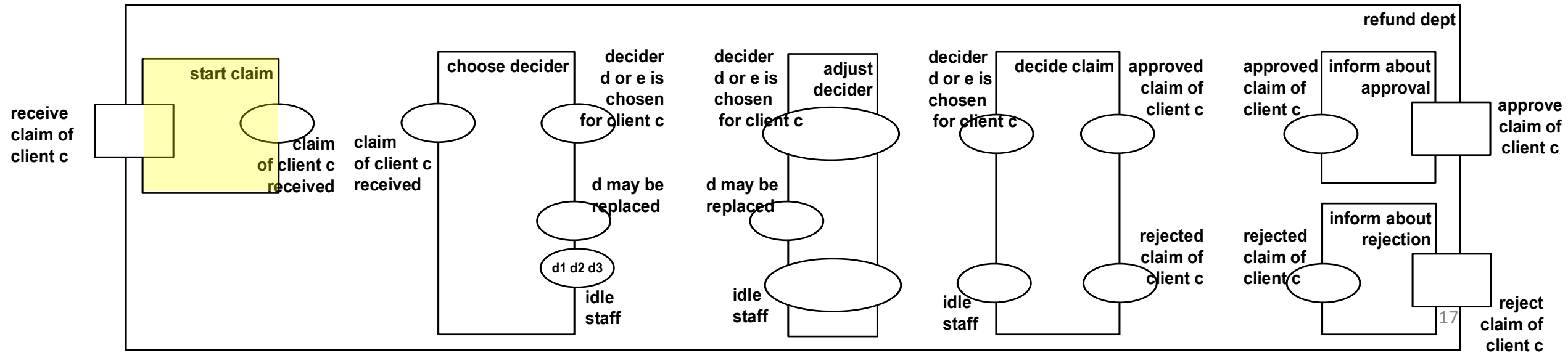
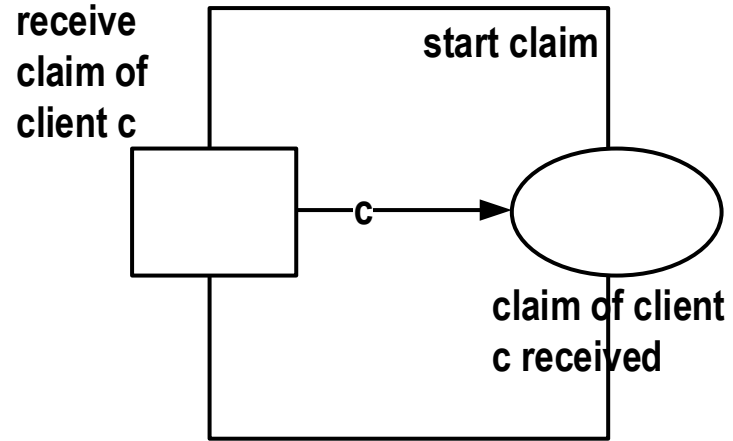
$$\text{il-runs}(M \bullet N) = \{a-b-c, b-a-c, a-c-b\}$$

$$\text{il-runs}(M) = \{a-b, b-a\} \quad \text{il-runs runs}(N) = \{c\}$$

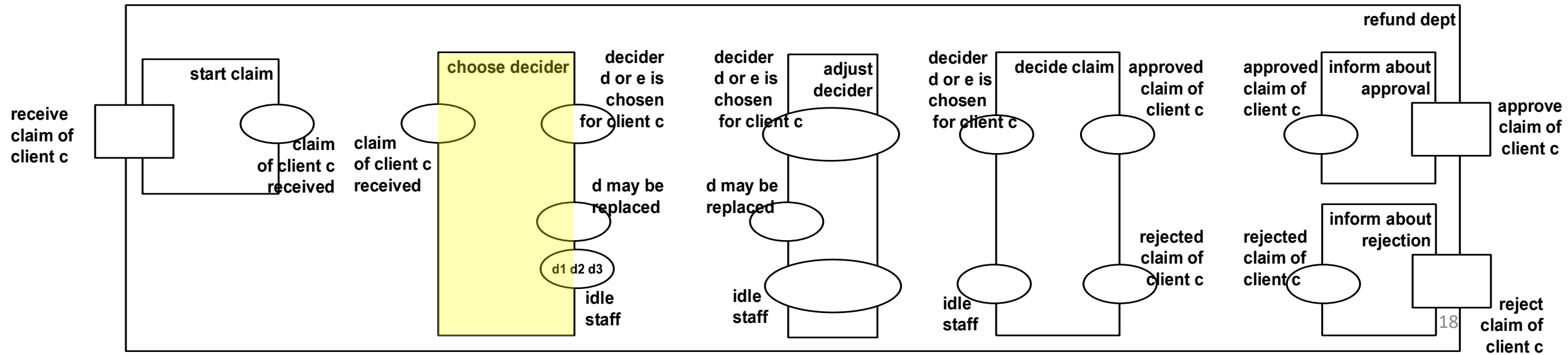
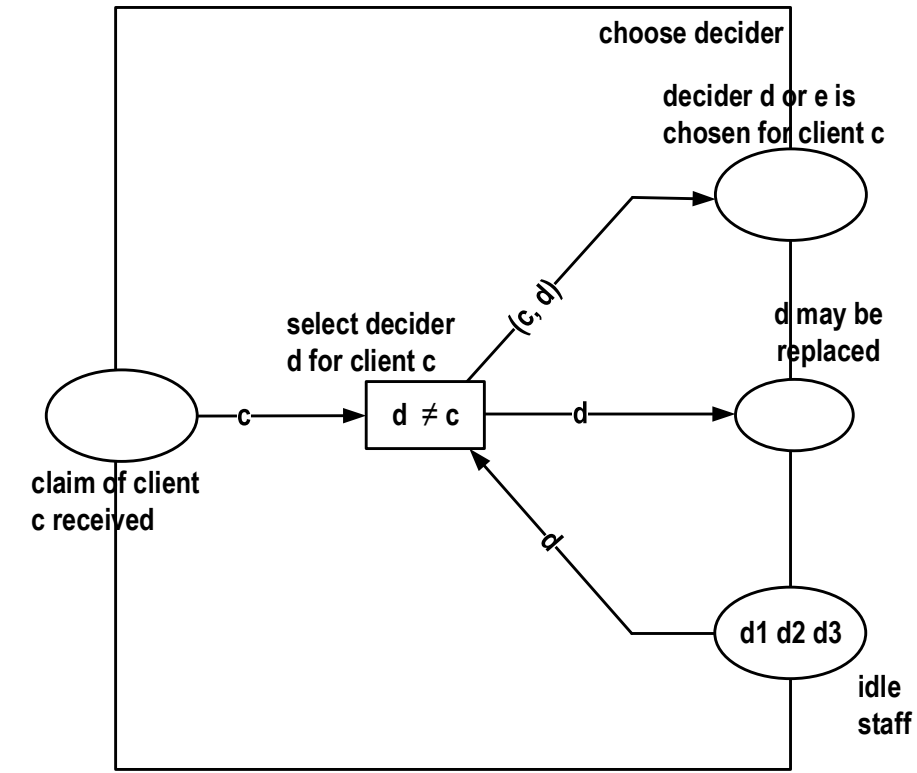
$$\text{il-runs}(M) \bullet \text{il-runs runs}(N) = \{a-b-c, b-a-c\}$$

$$\text{runs}(M \bullet N) = \text{runs}(M) \bullet \text{runs}(N)$$

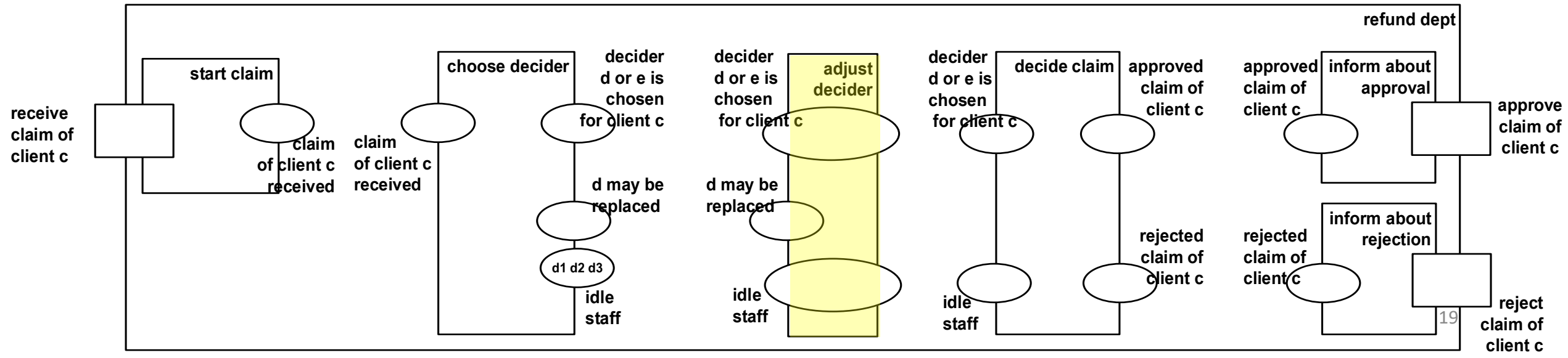
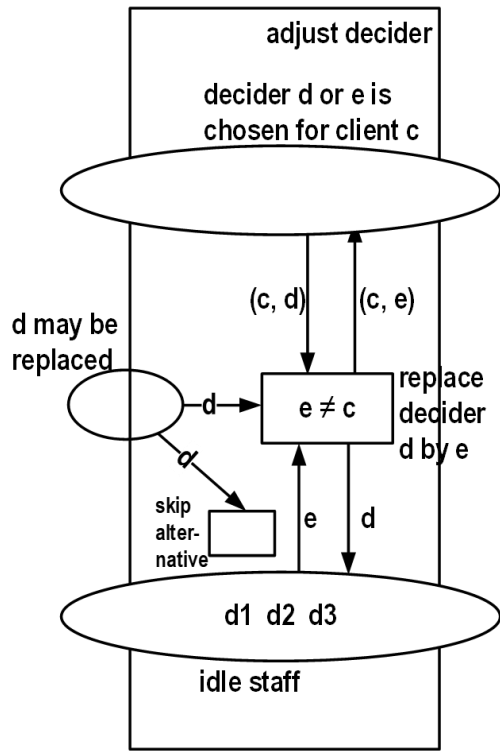
Part II system modules



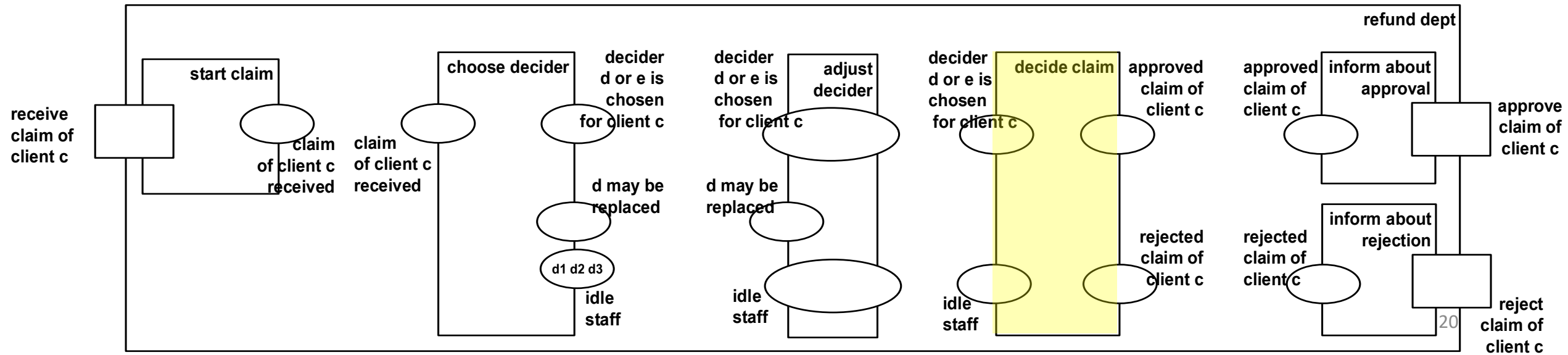
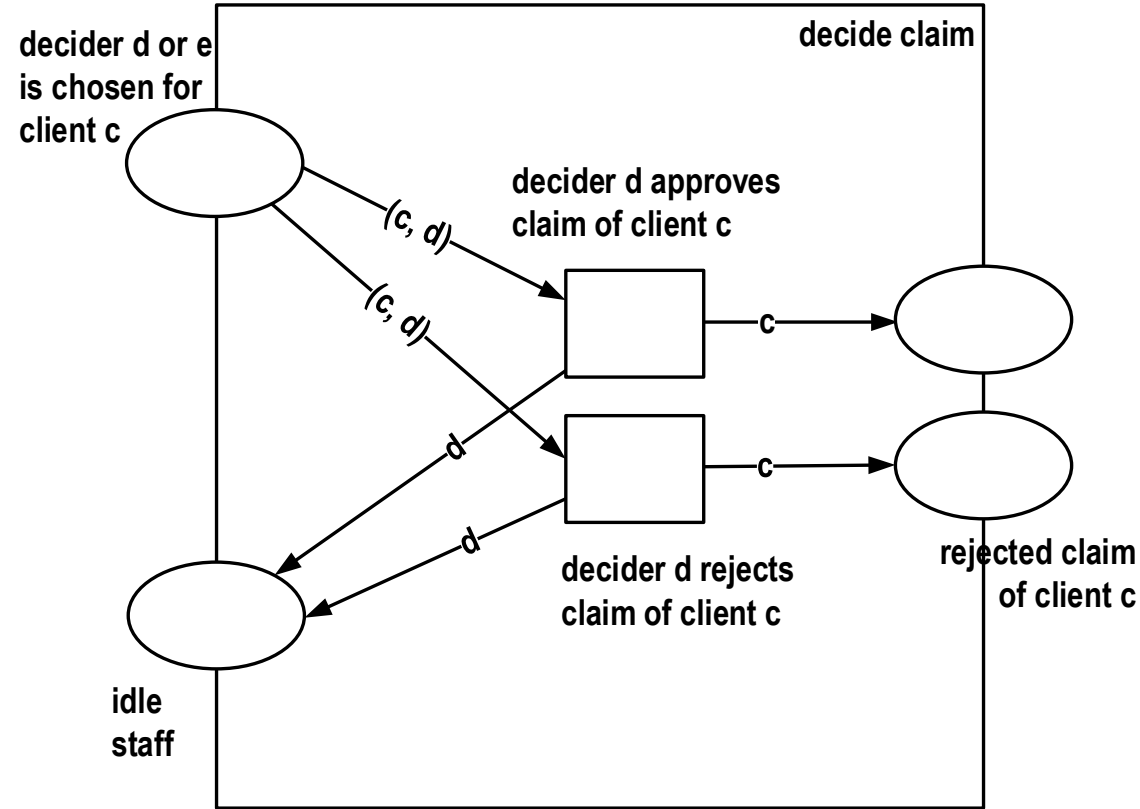
Part II system modules



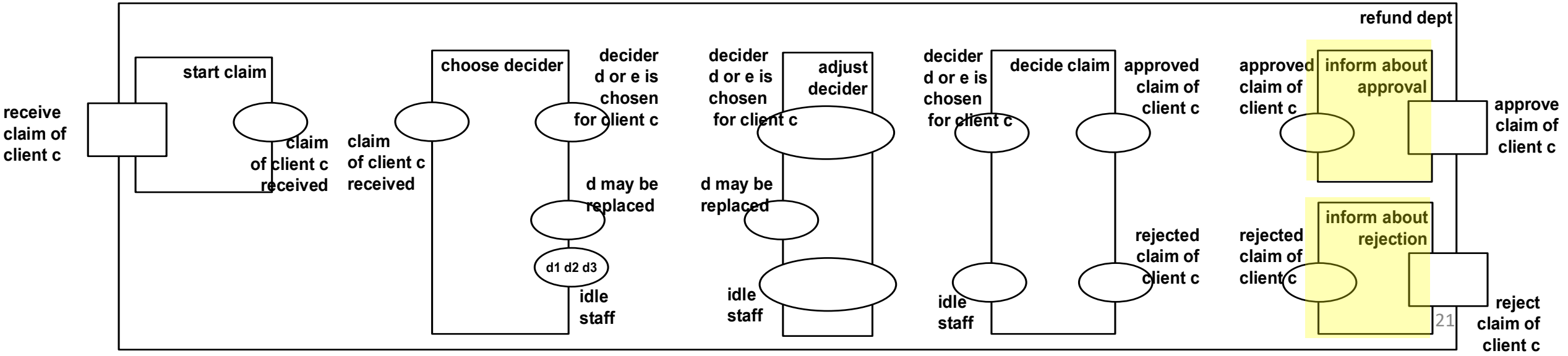
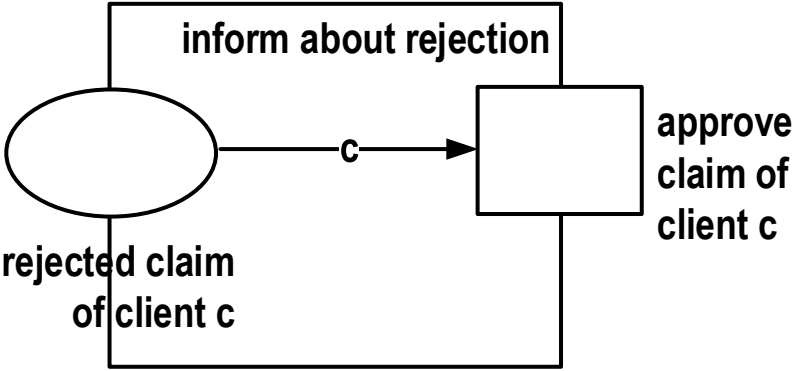
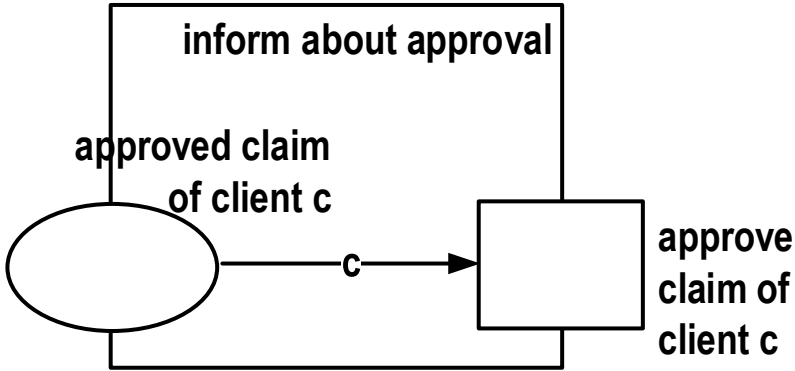
Part II system modules



Part II system modules

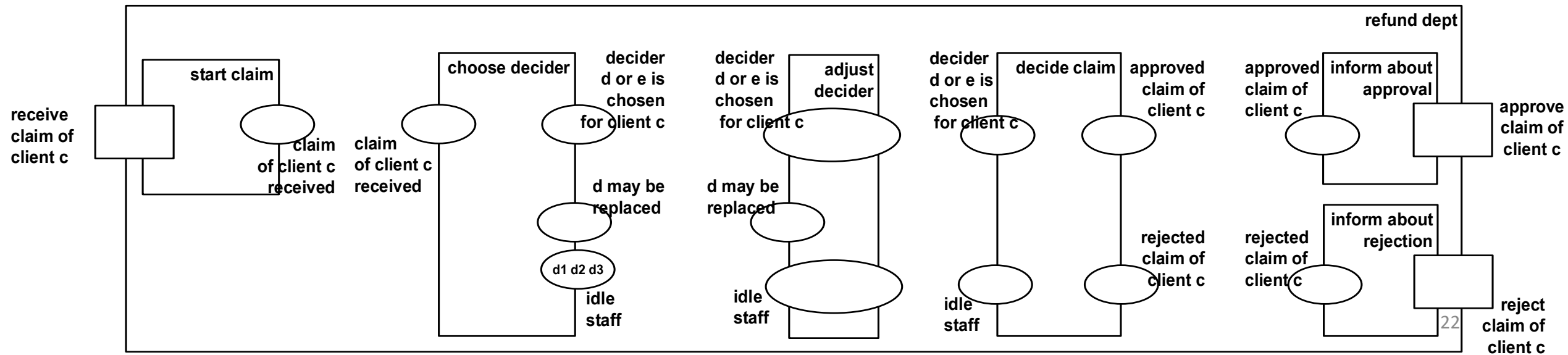


Part II system modules



Part II system modules

refund department =
start claim • choose decider • adjust decider • decide claim • inform about approval •
inform about rejection



Part III liveness

refund department $\models$

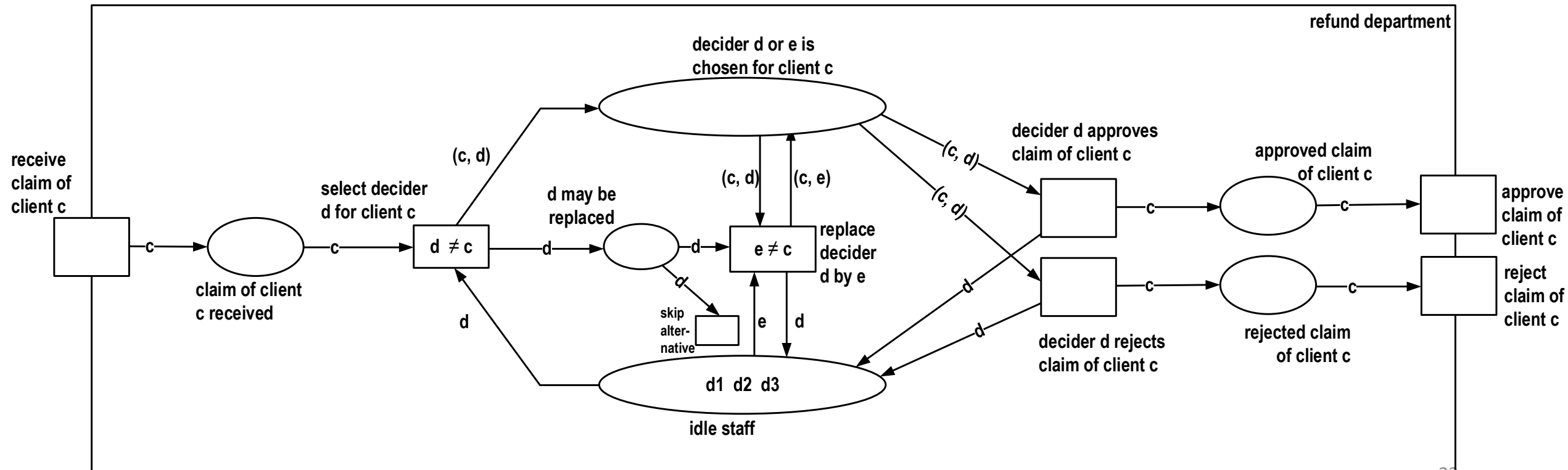
$\text{occurs}(\text{receive claim of client } c . c1) \mapsto$

$\text{occurs}(\text{approve claim of client } c . c1)$

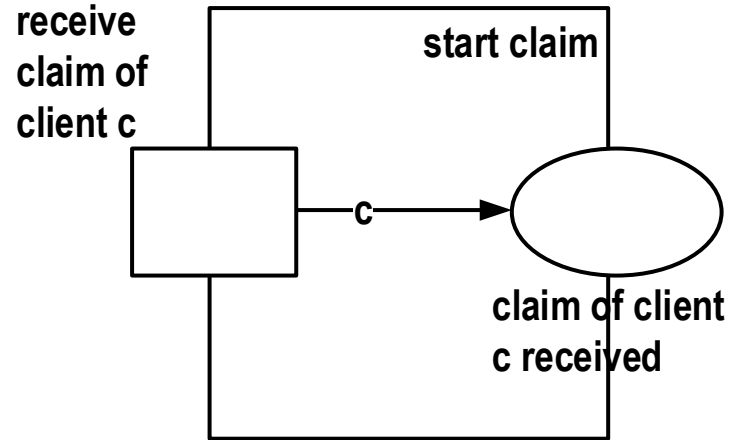
$\vee \text{occurs}(\text{reject claim of client } c . c1)$

How to prove this property?

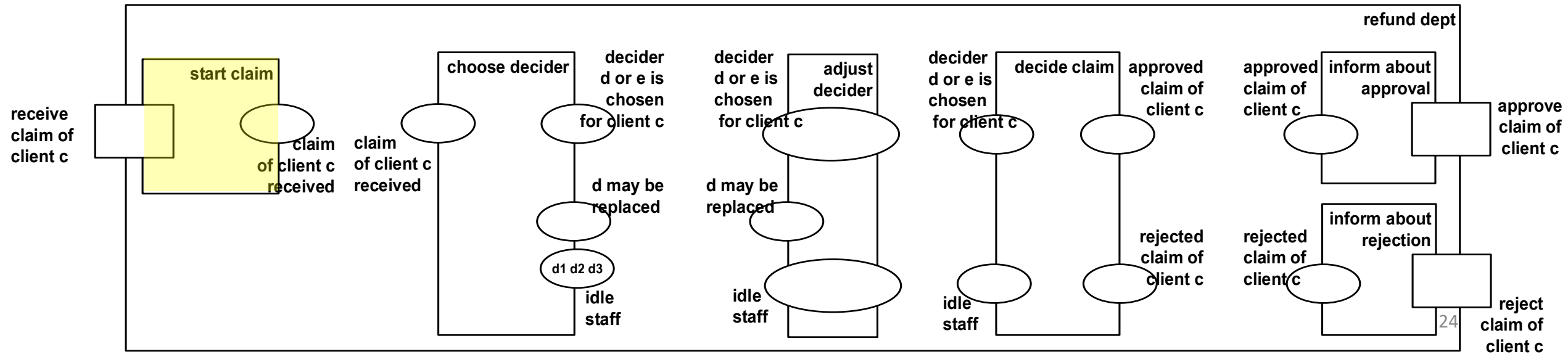
!Compositionally!



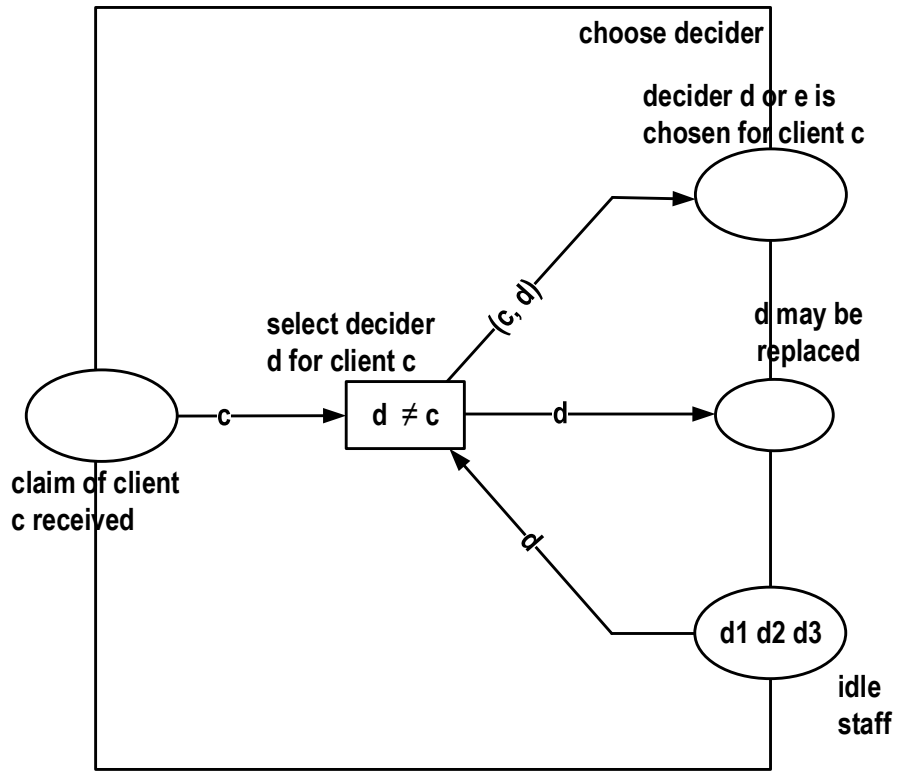
Part III liveness



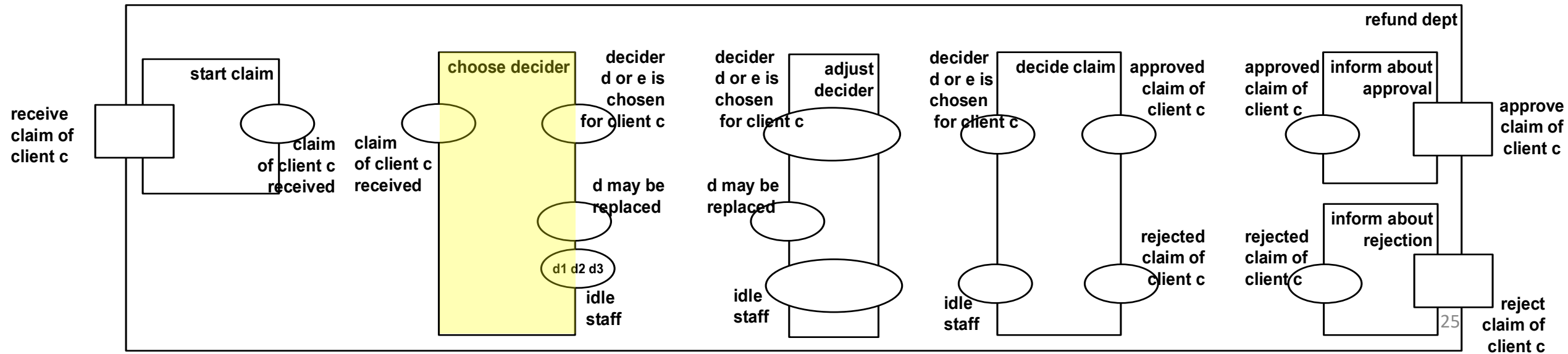
start claim $\models$
 occurs(receive claim of client c . c1)
 $\mapsto$ claim of client c received.c1



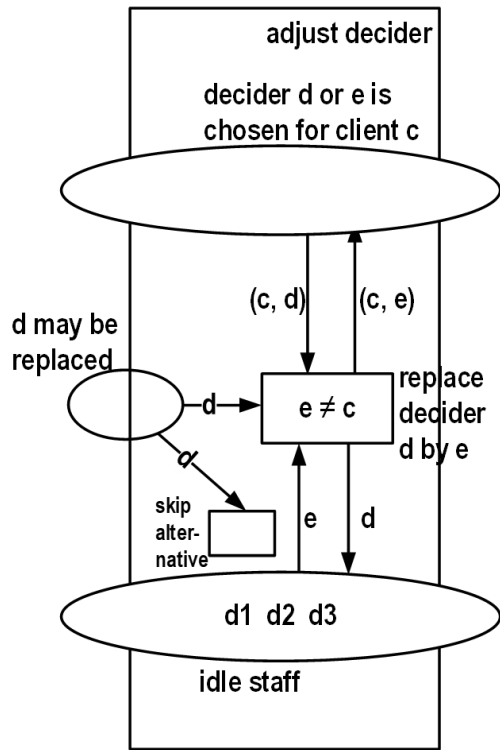
Part III liveness



choose decider $\models$
 idle staff.e $\wedge$ claim of client c received.c1
 $\mapsto$
 decider d or e is chosen for client c
 $\wedge$ d may be replaced



Part III Liveness



adjust decider $\models$

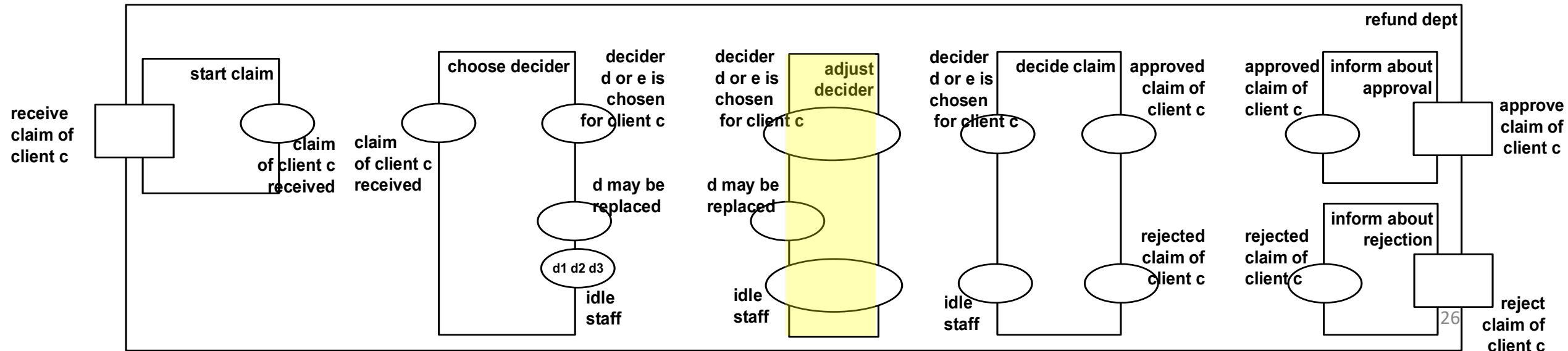
idle staff. d

$\wedge$ decider d or e is chosen for client c .d

$\wedge$ d may be replaced

$\Rightarrow$

decider d or e is chosen for client c. d



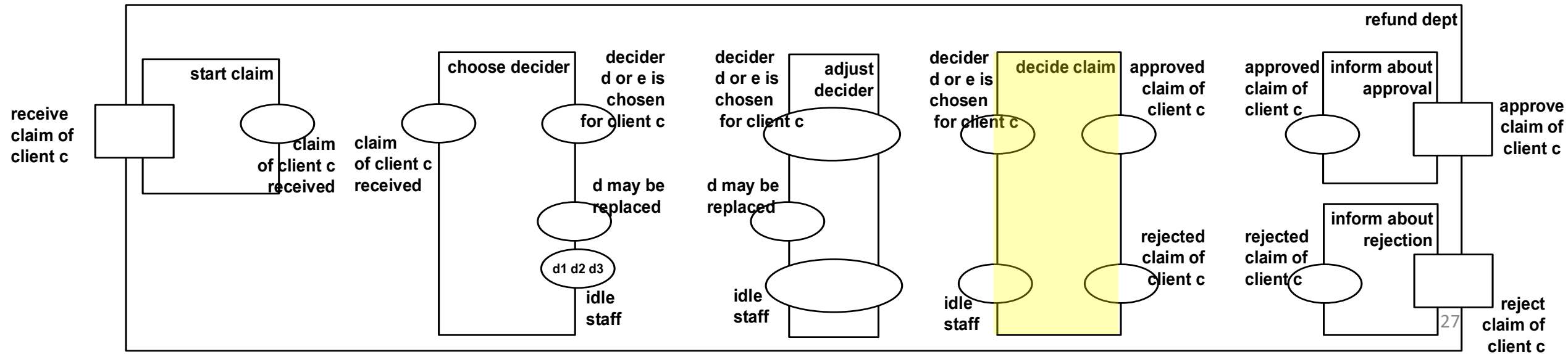
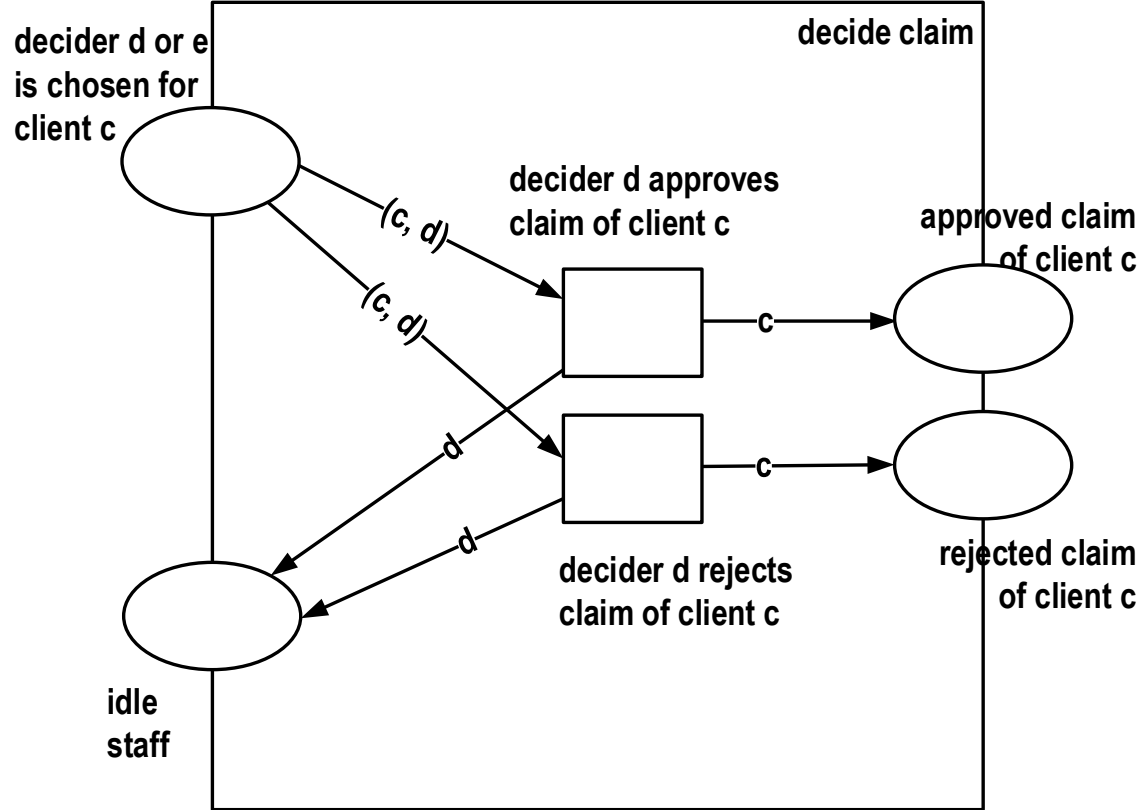
Part III liveness

decide claim $\models$

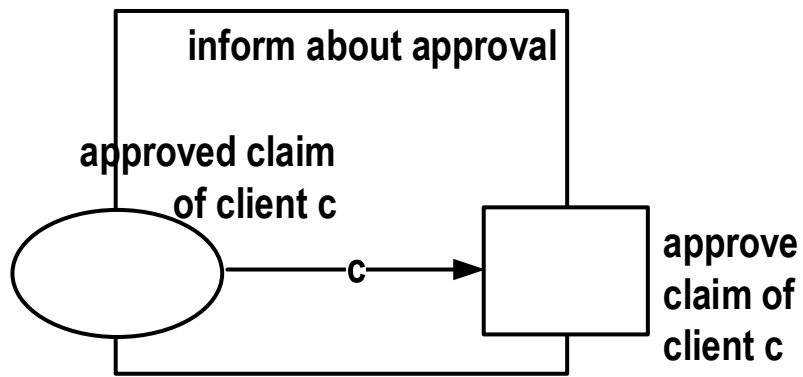
decider d or e is chosen for client c. d

$\mapsto$ approved claim of client c.c1

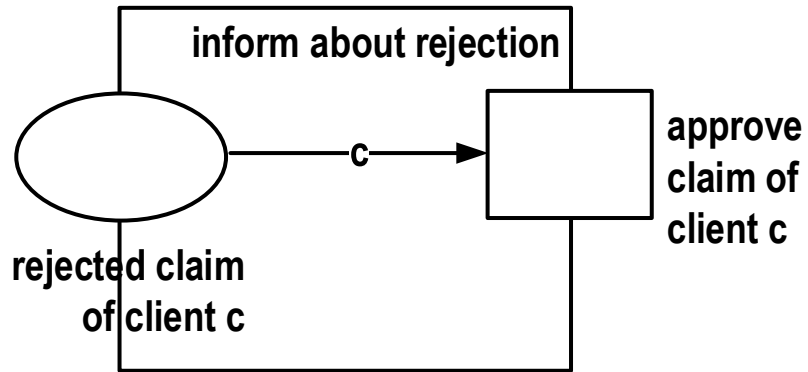
$\vee$ rejected claim of client c.c1



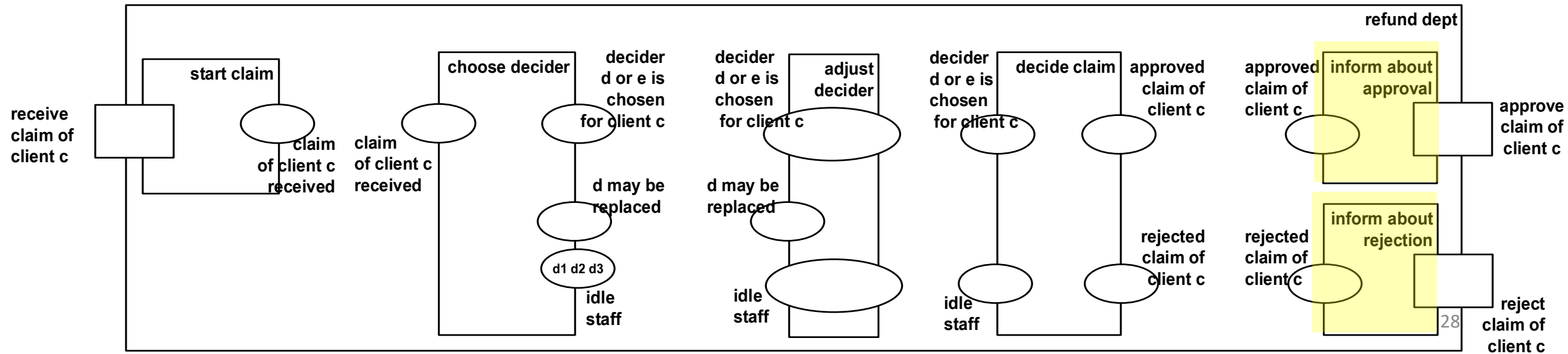
Part III liveness



inform about approval $\models$
 approved claim of client c . $c1$
 $\mapsto$ occurs(approve claim of client c . $c1$)



inform about rejection $\models$
 rejected claim of client c . $c1$
 $\mapsto$ occurs(reject claim of client c . $c1$)



Theorem Let M and N be two modules
with $M \models Mp \mapsto Mq$ and $N \models Np \mapsto Nq$.

Let $Mq \mapsto Np$.

Then $M \bullet N \models Mp \mapsto Nq$.

... *cum grano salis*

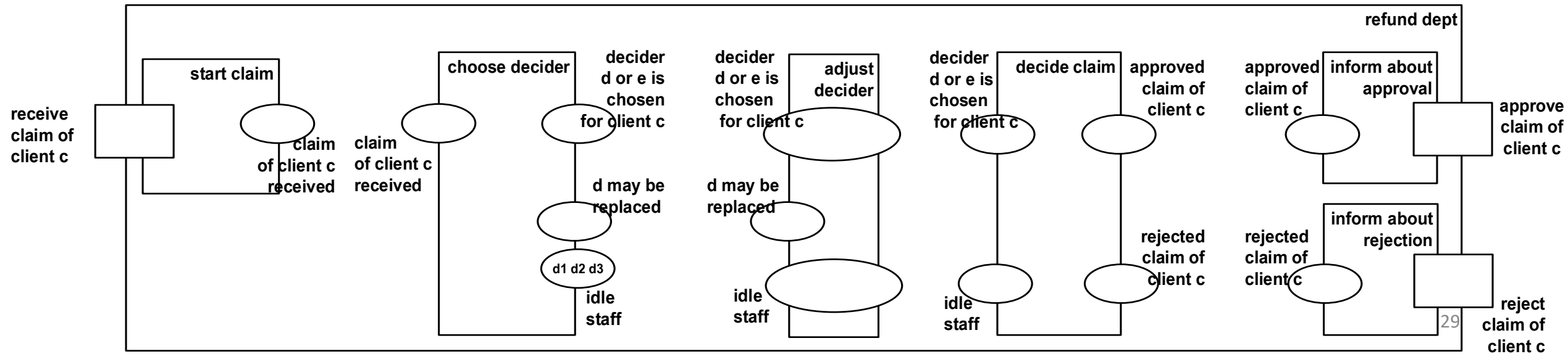
Part III liveness

refund department $\models$

occurs(receive claim of client c . $c1$) $\mapsto$

occurs (approve claim of client c . $c1$)

$\vee$ occurs(reject claim of client c . $c1$)



What did we gain?

- Modular Composition of modules

analogy:

words u, v, w over an alphabet Σ :

w can be composed at two ends: $u \bullet w$ and $w \bullet v$

Furthermore, $(u \bullet w) \bullet v = u \bullet (w \bullet v)$,

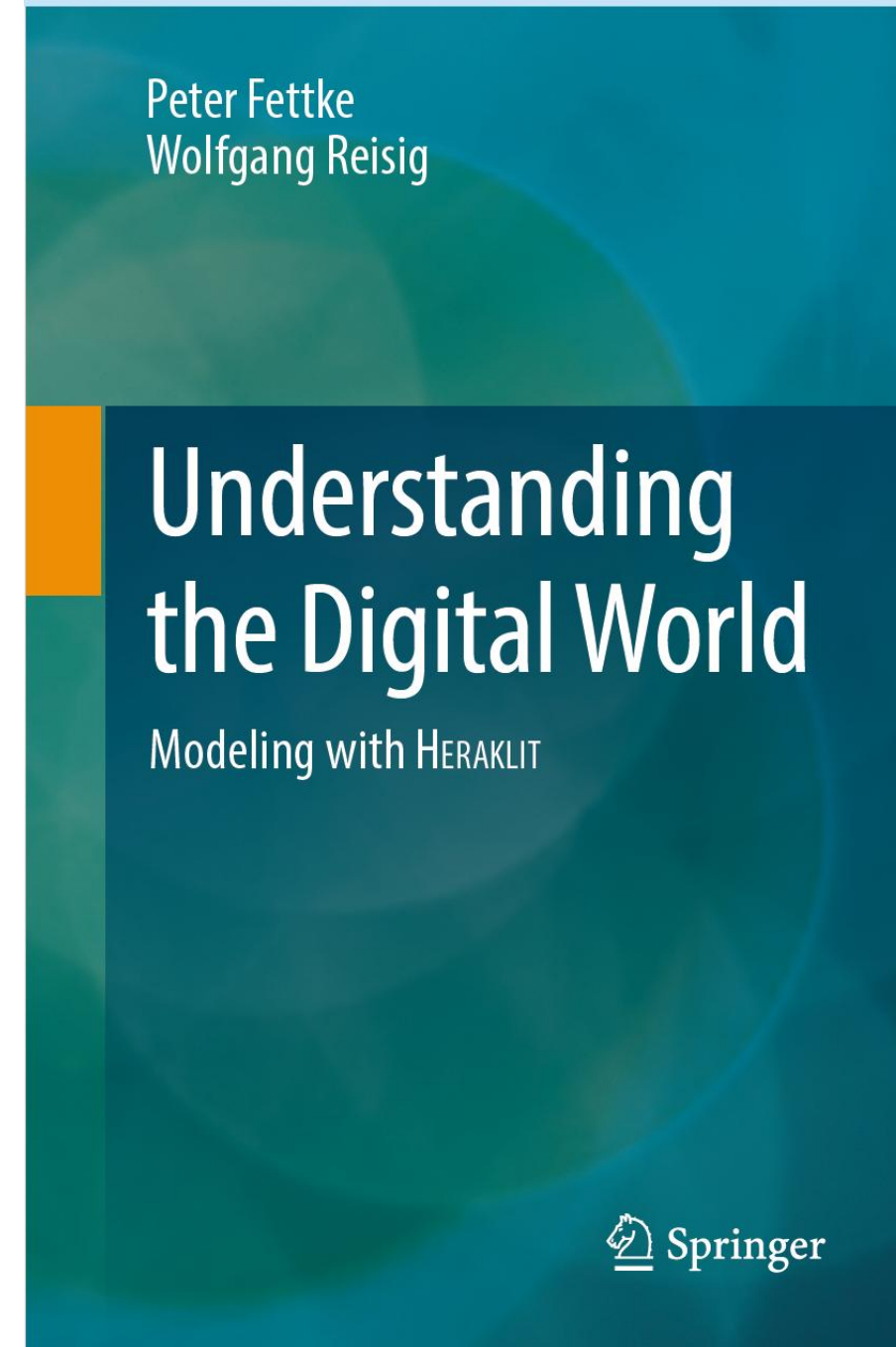
written $u \bullet w \bullet v$ or just uwv

Replace words with modules!

What did we gain?

- Modularization, Refinement, abstraction
- universal, expressive composition
total, deterministic, associative
- an algebraic representation of modules
- modular analysis techniques

for details, see



Trust by design

- in praise of modularization: a case study

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the end

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